

# COMMODORE 64 MACHINE LANGUAGE TUTORIAL

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A self tuition book on Machine Language Programming  
with the 65XX family of microprocessors.  
Suitable for use also with other Commodore, & Apple computers.

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# COMMODORE 64 MACHINE LANGUAGE TUTORIAL

BY PAUL BLAIR

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# COMMODORE 64 MACHINE LANGUAGE TUTORIAL

## PREFACE

**PAUL BLAIR** has distilled his own starting experiences in learning machine language (ML) into a set of TUTORIALS to ease the transition from BASIC to ML. The layout of the book is designed to provide self tuition and guidance in the simple use of the principal 6500 microprocessor family ML instructions. A computerist will gain confidence to continue into the language.

To assist discussion about ML programming, the book contains a complete program listing for the Commodore 64 (C64). There is a similar program available for the CBM/PET 4000 series of computers. As the Apple computer is also based on the 6500 family, Apple users can also learn from it. The point of the program is illustrative, and study of it will quickly demonstrate many of the points made in the text.

My contact with Paul Blair is as editor of the Commodore Magazine. After our initial introduction by mail, I was soon to find that Paul had a rare gift of being a natural born teacher - a gift he uses very well in this book.

**MERVYN BEAMISH**  
Editor Commodore Magazine

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# INTRODUCTION

My first toddling steps in machine language were disasters. Those curious listings that kept cropping up in the computer press were unintelligible, the magazines assumed I knew it all, and the books in the shops, mainly of the reference variety, gave little or no help to the novice. The rest is history now, but the further I went, the more conscious I became of the need for a text for the rank amateur. What follows in this book is a distillation of my starting experiences, set down to ease your transition to machine language (ML) programming.

The layout of the book is designed to provide guidance in the simple use of the principal 6500 microprocessor family ML instructions. You should gain sufficient confidence to go on from there.

To assist discussion about ML programming, the book contains a complete program for the Commodore 64 (C64). There is a similar program available for the CBM/PET 4000 series of computers. As the Apple computer is also based on the 6500 family, Apple users can also learn from it. The point of the program is illustrative, and study of it will quickly demonstrate many of the points made in the text.

Skill in ML programming can be achieved only by working at it - there is no silver bullet solution. Your level of competence will equal the effort you put into it.

PAUL BLAIR

# AFTER BASIC — What Next

By now, most of you will have become competent in the use of BASIC – and now wish to widen your programming skill horizons. BASIC is a high level language, and its designers have provided a large number of helpful (and forgiving) features in it. You can edit and run most BASIC programs in a friendly environment – even if your program is full of errors, you will get to run some of it, fixing up errors as you go along.

Not so machine language. By writing programs directly in the computer's own language, not a language interpreted for it, YOU become the definer of program logic – and if that logic is not correct, you may so confuse the computer that you lose control of it. So, be warned, ML programming is demanding, exacting and often exasperating. But it is a great teacher, too. Your logic processes will be strengthened, your temper becomes more elastic and your self satisfaction will grow with each successful excursion.

Why ML? What can you manage with ML that can't be done in BASIC? The answer to the latter question is – nothing! In fact, BASIC provides many routines that you would have to provide yourself in ML. BASIC will multiply 2 numbers for you by saying

```
'PRINT 2*3'
```

whereas ML requires you to design a multiply routine.

The advantage of ML is principally speed. By removing the intermediate interpretation step of BASIC, your computer can run hundreds, if not thousands of times faster. For tasks like sorting, searching and some graphics, ML is the fastest (and probably the only satisfactory) way of programming. All tasks that require a large number of repetitions of whatever type – that is where ML will win hands down. You will also find that judicious use of ML permits you to use larger or more complex programs.

You need a few items, so that you can experiment for yourself. The first is easy – get hold of a small hexadecimal to decimal conversion chart, one that can sit just above the keys on the keyboard (or be taped there). The other is an extended monitor program, like 64MONITOR or SUPERMON. This will make learning easier.

There are some things that need to be covered before going too much further, so let's do them now.

# TOPIC 1 – Memory Organisation

We are going to concentrate on the '65XX' family of microprocessors. The original Commodore computers (CBM/PET) and the VIC used the 6502 (as did the Apple computer). The Commodore 64 uses the newer 6510 microprocessor. For our purposes, the 6502 and 6510 are the same, and '65XX' as used here will refer to both devices.

Central to an understanding of ML is some idea of computer design, particularly how information is stored.

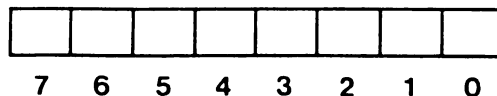
The original storage concept has not varied from the earliest days of data processing, when small magnets were 'core' storage. If the magnetic field went one way, the magnet was 'on', and the other way was 'off'. This idea persists, only the storage medium has changed.

The modern computer relies on the 'on/off' principle. Each 'bit' (Binary digit) or storage location can be 'on' or 'off'. By convention, 'on' is as if the bit represented one, and 'off' as if the bit represented zero.

For ease of processing, bits are arranged in groups, named 'bytes'. In the 65XX, eight bits makes a byte – so Commodore computers are referred to as '8 bit' machines. Later or larger computers use 16 or 32 bits in a byte.

So let us represent a byte like this. Each small box represents a bit, and the large overall box represents a byte.

Figure 1



For convenience, each bit is given a 'position number', starting with 0 on the right running to 7 on the left (this is akin to decimal units, tens, hundreds and so on). But there is a reason, as we shall see.

How do we use this? As with decimal numbers, each bit (whether on or off) has a 'weight'. In fact, the bit position numbers represent these weights. Because each bit can only represent 2 conditions, on or off (or 1 and 0) the multipliers increase in powers of 2. In this way, the value of a byte is formed.

Using the bit position numbers, the individual bits either represent 2 raised to the power of the bit position, or zero. So Bit 0 can represent 1 or 0, bit 1 can mean 2 (2 to the power of 1) or 0. Bit 2 can mean 4 (2 to the power of 2) or 0, bit 3 can mean 8 or 0, and so on. Bit 7 can mean two to the power of 7 ( 128) or 0. So we have a range of number combinations in each byte, depending on which bits are on or off. If all possible bit combinations of 'on' and/or 'off' are considered, and the values of individual bits are added together, each byte can mean any number from 0 to 255 (try it yourself and see what you get). To put it more graphically, let's do some examples.

### TASK 1

Draw bit patterns for the following numbers-

200 2 89 0 254 32 176

What numbers do these bit patterns represent-

11000101  
00011011  
10101010  
01010101  
11110000  
00001001  
10010010

Decimal numbers are one thing to remember, but who would want to remember bit pattern numbers? There had to be an easier and tidier way.

One byte can represent 0-255, i.e., 256 numbers. 256 is 16 x 16 ..... so every number in the range 0-255 could be represented by not more than two 'digits' if we could think of a way to represent a range of 0-16 with each 'digit'.

# TOPIC 2 — Hexadecimal Notation

We have been conditioned to think in numbers of base 10 – i.e., numbers in ‘groups’ of tens. Any base could be used, but the world has used base ten for normal mathematics. Why base ten and not base 12 or 7 probably stems from the number of fingers we have! But computers don’t have fingers, just bits!!

The 16 x 16 product led to thoughts of base 16 – but there weren’t enough numerals to go around. To overcome this, the letters A to F were ‘added’ to the range of numerals. As if counting in 16’s wasn’t enough!

Instead of counting in tens, we count in 16’s. Consider the range of decimal numbers from 0 to 15. Let us draw a table showing equivalent hexadecimal (usually shortened to ‘hex’) numbers:-

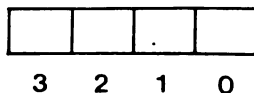
DEC	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
HEX	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F

As with decimal, we add another position (‘tens’ in decimal) when we go past the limit of the base number. This idea extends to hundreds, thousands, and so on in decimal, so why not hexadecimal? To count on from where we left off-

DEC	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
HEX	10	11	12	13	14	15	16	17	18	19	1A	1B	1C	1D	1E	1F

We can draw a comparison with our description of the bit ‘weight’ in a byte. Just as each position has a weight or indice, starting at 0 and incrementing by 1 for each step to the left, so hexadecimal numbers have weights for a similar pattern, but this time the ‘weight’ is in 16’s. Figure 2 shows how.

Figure 2



Just as in our example with the bits in a byte, so hexadecimal applies 'weights' in relation to the position number. Think of the position number as representing the lowest value the position can hold – so position 0 starts from 0, position 1 starts from 16, position 2 starts from 256 (16 x 16) and position 3 from 4096 (16 x 16 x 16).

To convert from decimal to hexadecimal, divide the decimal number by 4096, 256 or 16, whichever is the next smallest number to the number to be converted. The whole part of the result converts to its hex equivalent. The process continues with the next smallest divisor, until the final remainder is less than 16. The concept is simple, and with practice you will become familiar with its use. Let us convert a number, - say 5127

$5127 / 4096 = 1$  and 1031 over = 1 as first hex digit  
 $1031 / 256 = 4$  and 7 over = 4 as next digit  
 $7 / 16 = 0$  and 7 over = 0 as next digit

So the final conversion of 5127 is 1407 hexadecimal.

The reverse applies – multiply progressively by 16, 256, and (if needed) 4096, adding together the individual results.

There are certain conventions used in hexadecimal notation. To differentiate between decimal and hexadecimal, hex numbers are prefixed with '\$'. And hex numbers are usually written in groups of 2 numerals – \$XX or \$YYYY, not \$Z or \$ZZZ. This is not an absolute rule, but a tidy one.

## **TASK 2**

Convert some numbers to hex–

10 25 49 8 100 47 19 66

Convert some hex numbers to decimal–

\$4E \$07 \$5B \$6A \$49 \$17 \$25

You will come to grips with hexadecimal with practice. No book will do this for you, it's a matter for personal effort. Don't try to remember it all – hence the earlier suggestion for a small hex/decimal table to tape onto your computer keyboard.

# TOPIC 3 — The 65XX Microprocessor

You probably have looked into the heart of the main processor chip, either accidentally or deliberately. To do this, you would have accessed the monitor, built into 3000/4000/8000 Series Commodore machines, or available as a loadable program for the C64. The monitor is the window on the 65XX, the chip that is the heart of Commodore computers. If you are to succeed in ML programming, you need some basic information to start – and you will learn a lot more as we go along.

When you got to the monitor and peered in, you probably got a picture something like this

```
PC  IRQ  SR  AC  XR  YR  SP
DC4D EA31 CC 30 05 FF F8
```

and (no doubt!) wondered what you had struck. You weren't the first, and won't be the last, let me assure you.

Storage in the 65XX can be considered as a number of memory locations, hooked together by a common line (the data bus) on which all information travels. In any program, the 65XX (by its 'architecture' or the way its brain is arranged) uses the 3 main temporary mailboxes AC, XR and YR to hold information, counting where it's at with the PC, checking its filing system with the SP, watching the instruments with the SR, and keeping overall control with the IRQ. Understand all that? Thought not.

Let's go through all that again, with some explanations.

*PC stands for Program Counter.* Using this counter, the computer stores enough information (2 bytes, each of 8 bits) to remind itself what to do when it finishes what it is doing now. In the example, \$DC4D is the next port of call.

*IRQ is short for Interrupt ReQuest.* When a program is running, it is interrupted momentarily every so often (like 60 times a second) to check for outside activities, such as whether someone has typed something at the keyboard. In that case, whatever has been typed needs to be stored, then the program will continue. Don't worry too much about it for now, just be aware that it exists, and is very important.

*SR stands for Status Register,* which is one byte whose bits can be set or unset ('on' or 'off' if you like) to denote certain things – e.g., if bit 7 is 'on' (= 1) it could denote that the number you are working on is negative, while bit 7 'off' (= 0) could indicate a positive number. The bits in the SR are termed 'flags', which we will discuss in Topic 10.

The next 3 locations are also named registers, and are one byte each.

*AC is the ACcumulator,* probably the most important register of the three. For convenience, the name is abbreviated to '.A' During programs, most operations centre on .A

*XR and YR are the X and Y Registers,* usually referred to as .X and .Y Like .A, they are each one byte of storage, used for holding variables during run time.

*SP stands for Stack Pointer.* The stack is a 256 byte block of storage used for holding all sorts of information. It works like a pile of dinner plates – the last plate you put on top will be the first plate you retrieve. Information is 'pushed' onto the stack one byte at a time, and retrieved (or 'pulled') the same way, but in the reverse order to the order in which it was stored. Sometimes called LIFO, or Last In First Out, which sums it up pretty well. Of course, this is not to be confused with GIGO (Garbage in Garbage out) – a well known computer disease.

That completes our introductory sections. Reread these notes from time to time, to keep them fresh in your mind.

# TOPIC 4

## — Machine Language Programs

Commodore computers were designed around the 6502 (more recently the 6510) chip, and so enjoy the use of what is termed the '6502 instruction set'. The 6510 and the 6502 are identical in this respect. In brief, this means that it is possible to instruct the computer to perform certain functions, such as add, subtract, move data and so on, by using an 'instruction' that is included in the computers vocabulary. The instruction set has been carefully designed to provide for most functions that are needed, but (as ever, I suppose) there are a few not provided for directly, but are available by combining other directions. This is not a serious limitation.

The instruction set contains 56 core instructions, and a number of these have subsets of operating parameters, so that the 56 instructions have a total of 151 different operating techniques. That probably sounds daunting, but there is a pattern in the subsets that makes it much easier to follow. Some texts use the subsets to introduce the core instructions, but we will work the other way, using the instructions to introduce the subsets.

The 56 core instructions are defined by 3 character alphabetic codes, named operating codes – opcodes for short. In the sections that follow, we will examine each major opcode in turn.

### *THE OPCODE SET WITH TRANSLATIONS*

× ADC	ADd memory to .A with Carry
AND	Logical AND of memory with .A
ASL	Arithmetic Shift memory of .A Left one bit
– BCC	Branch if the C (carry) flag is Clear (0)
– BCS	Branch if the C flag is Set (1)
– BEQ	Branch if the Z (zero) flag is EQUAL to 1
BIT	Test memory bits
– BMI	Branch if the N (negative or MInus) flag is 1
– BNE	Branch if the Z flag is 0 (result Not Equal to 0)
– BPL	Branch if the N (negative) flag is 0 (result PPlus)

- BRK Force a BReak (stop)
- BVC Branch if the V (oVerflow) flag is Clear (0)
- BVS Branch if the V (oVerflow) flag is Set (1)
- CLC CLear the C (carry) flag to 0
- CLD CLear the D (decimal) flag to 0
- CLI CLear the I (interrupt) flag to 0
- CLV CLear the V (overflow) flag to 0
- CMP CoMPare memory with .A
- CPX ComPare memory with .X
- CPY ComPare memory with .Y
- DEC DECrement memory contents by 1
- DEX DEcrement .X by 1
- DEY DEcrement .Y by 1
- EOR Arithmetic Exclusive OR of .A and memory
- INC INCrement memory contents by 1
- INX INcrement .X by 1
- INY INcrement .Y by 1
- JMP JuMP to new location
- JSR Jump to new location, Saving the Return address
- LDA LoAD .A with a byte
- LDX LoAD .X with a byte
- LDY LoAD .Y with a byte
- LSR Logical Shift memory or .A Right one bit
- NOP No OPeration – do nothing
- ORA Inclusive OR memory with .A
- PHA PusH value of .A onto stack
- PHP PusH value of SR (Processor) onto stack
- PLA PuIL top value off stack into .A
- PLP PuIL top value off stack into SR (Processor)
- ROL ROtate memory or .A one place to the Left
- ROR ROtate memory or .A one place to the Right
- RTI ReTurn from Interrupt
- RTS ReTurn from Subroutine
- SBC SuBtract memory from .A with Carry
- SEC SEt the C (carry) flag to 1
- SED SEt the D (decimal) flag to 1
- SEI SEt the I (interrupt) flag to 1
- STA STore the value in .A into memory
- STX STore the value in .X into memory
- STY STore the value in .Y into memory
- TAX Transfer .A to .X

- TAY      Transfer .A to .Y
- TSX      Transfer SP into .X
- TXA      Transfer .X into .A
- TXS      Transfer .X into SP
- TYA      Transfer .Y into .A

Most opcodes have to operate on something – the opcode is not complete in itself, with a few exceptions. Rather than set it all out here, the ideas of labels and operands (as the other parts of a typical instruction are termed) will be set out as we go through the opcode family.

The list looks rather unnerving, so without further ado, let's set about demolishing it.



# TOPIC 5 — Load and Store

## LOAD

The load and the store operations are very similar, with one important (and obvious) difference.

You can load .A, .X and .Y with either a value, e.g., 20, or you can load the actual value that is already stored in another memory location.

If you want to load a specific value, say decimal 32 (\$20) you would write

```
LDA # $20 (Type 1)
or LDX # $20
or LDY # $20
```

This introduces the '#' (hash) symbol. Put simply, it differentiates between the address of a location (\$20) and the value hex \$20 (#\$20). I prefer to work in hex (then there is no mistake about whether I'm in decimal or hex notation), so I will continue to use the hash convention here.

It also enables an introduction to the term 'operand' – which is simply something to be operated on. The '#\$20' is operated on by LDA – and whether it is a value or a memory location, it carries the same title.

So, if the status value (ST) that the C64 uses to detect the end of a file is stored in memory location \$90 (\$96 in CBM/PET), you can call that value into any of the 3 registers with

```
OPCODE.OPERAND
  LDA $90      (Type 2a)
  or LDX $90
  or LDY $90
```

Of course, if you can load from a memory location with a 2 hex digit address, you can also load from one with a 4 hex digit address – giving

#### OPCODE.OPERAND

LDA \$B000 (Type 2b)  
or LDX \$B000  
or LDY \$B000

Each type of operation conforms to a specific type of operation. Not all texts use the same naming convention, so I will use the most common one.

Type 1 statements are uniquely named 'immediate' i.e., the load a value specified in the statement – Load Accumulator with the number represented by \$20, or 32 decimal. Very straightforward.

Type 2 statements are named 'absolute' – what the statement says it is! Go to that location, get the value found there, and put it into the correct register. Again, no problem.

Remember that each byte can hold only the range of values from 0 to 255.

The next form of load opcode is probably one of the most useful. In this form, an index is used. The index relates to the address specified in the instruction, and is added to it. To be more specific, consider the following –

LDA \$B000,Y (Type 3)

Depending on the value in the .Y register, load .A with the value found by adding .Y to \$B000. If .Y held 7 (#\$07 by our convention), then the instruction becomes –

LDA \$B007

Note that by performing the addition, I can remove the comma, turning a Type 3 statement into a Type 2 statement. But if I can somehow manage to arrange a loop in which I increase .Y, then I can sequentially load consecutive bytes into .A

Type 3 instructions are termed 'indexed absolute' and you can use .X and .Y (but NOT .A) as indices. You wouldn't use LDX \$B045,X-it wouldn't work.

The final form in this group is really only a change of name. The 256 bytes from \$00 to \$FF (at the very start of computer memory) are called 'zero page'. Usually the locations in zero page are used by BASIC programs to hold all sorts of useful information – the status word ST is one. If BASIC is not being used, many of the locations are sitting idle, and are an ideal place to store useful values in a ML program. So, using zero page is useful, and leads to Type 4 addressing (which looks rather like Type 3)

LDA \$A0,X

This is termed 'indexed zero page' – in reality 'indexed absolute' using zero page. Zero page addressing is useful because it uses 1 less byte in each instruction, and operates more speedily than a reference to any other part of memory.

## STORE

As noted earlier, store works like load, with one difference. While you can load any number or from any memory location you wish, (immediate type) you obviously cannot do an identical reverse, such as writing to fixed ROM. But you can store a variable into any random access memory location. Consider the following –

STA \$07  
STY \$0401  
STX \$E007

The first two will work – in the C64 location \$07 is in RAM, and \$0401 is the second position in screen RAM. But \$E007 is Read Only Memory, and you can't normally write to ROM, so the instruction will fail.

Just as you can perform an indexed load, so you do an indexed store –

STA \$0401,X

is quite legal, as is

STX \$07,Y

### SUMMING UP

You have seen lots of LDA's, STX's and so on, and you might (at this time) think that the same rules apply to all instructions. The fact is, they don't. Not every instruction supports the same addressing modes as its brothers, so let's set out a table for future reference. The table will list each command, with 'Y' or 'N' to indicate whether the mode is legal or not.

OPCODE	IMMED	ABSOLUTE	ZERO PAGE	INDEXED ABSOLUTE		INDEXED ZERO PAGE	
				X	Y	X	Y
LDA	Y	Y	Y	Y	Y	Y	N
LDX	Y	Y	Y	N	Y	N	Y
LDY	Y	Y	Y	Y	N	Y	N
STA	N	Y	Y	Y	Y	Y	N
STX	N	Y	Y	N	N	N	Y
STY	N	Y	Y	N	N	Y	N

This makes it quite clear which goes with what. Don't try to remember them all – you can look up books or make up a collection from these notes for reference.

# TOPIC 6 — Increments and Decrements

The next 4 opcodes are termed 'implied', i.e., they have no operand following the opcode, as the operand is implied in the instruction. They are stand-alone instructions. The 4 instructions are –

INX      which adds one to .X  
INY      which adds one to .Y  
DEX      which subtracts one from .X  
DEY      which subtracts one from .Y

There are 2 other opcodes in this group. Like load and store, they have absolute, indexed and zero page options, and all operate on memory locations in RAM

INC \$033C adds one to location \$033C  
DEC \$A8 subtracts one from location \$A8

INC and DEC support absolute and zero page, and both of these instructions indexed via the .X register.



# TOPIC 7 — Register Transfers

The 65XX can shift values between registers with 6 'implied' instructions.

TAX	transfers .A to .X
TAY	transfers .A to .Y
TXA	transfers .X to .A
TYA	transfers .Y to .A
TSX	transfers SP to .X
TXS	transfers .X to SP

There is no direct transfer available between .X and .Y or vice versa. Note that in each transfer, the contents of the source register are transferred to the target register **WITHOUT** altering the contents of the source register.

There is no INA (INcrement .A) instruction, so transfer commands are combined to provide for this –

```
TAX ;MOVE .A TO .X
INX ;ADD 1
TXA ;AND RETURN IT TO .A
```

This, of course, loses the value in .X at the start. So, if .X is important at the time, you could –

```
STX $D7 ;STORE HANDY
TAX
INX
TXA
LDX $D7 ;GET IT BACK
```

INA would be a useful command – in this example it would save at least 7 bytes if used in a program. Later on, a substitute method will be shown.

## *SUMMING UP*

It might surprise you to know that of the 153 instructions we started out with, 41 types have been covered in this first round. If you don't believe me, count 'em !

# TOPIC 8 — Jumps and Branches

BASIC has some very useful commands that permit movement to other parts of programs— e.g., GOSUB takes you to special or often used routines, and RETURN gets you back home again. GOTO also takes you to other parts, straight off. IF..THEN or ON..GOTO permit you to apply some sort of test or check before going to some special place or other. Luckily for ML programmers, there are similar provisions in the 65XX. Let's explore further.

## JUMPS

The most commonly used jump instruction is JSR, which is Jump and Save Return. The ML program must remember the address that it jumped from, so that it can return there later. Sometimes JSR is described as Jump to SubRoutine, but that isn't the best description. However, just as a GOSUB works, so JSR takes you to those places in your program where you want particular things to happen, and brings you home again.

The brother instruction to JSR is RTS, which is ReTurn from Subroutine. RTS can be used to finish a subroutine and go back to the mainstream of a program, or, if all work is finished, to return to BASIC.

The last of the jump-type instructions is JMP, which is a straight GOTO. When you say JMP you mean just that— there is no RTS from a JMP. You can jump anywhere in your program, or jump to a routine in the machine itself, such as a routine to print a string on the screen.

## BRANCHES

We will treat branches at face value. In fact, they do more than I'm about to tell you, but we'll come back to that later.

Branches allow you to test some condition, and depending on the result of that test, the program will transfer operations backwards or forwards to another location. The only limitation here is that you can only go forwards or back to the limit of about half the maximum value that can be stored in a byte— so we can go back or forward about 127 bytes (near enough for now). This means that the point to which you want to branch has to be reasonably close by – you can't branch from one end of a long program to another.

⊗ NB DO NOT NEED TO CMP AFTER INX OR DEX  
The principal mathematical branch opcodes are—

BNE	Branch if the result is Not Equal to zero
BEQ	Branch if the result is Equal to zero
BMI	Branch if the result is Minus (< 0)
BPL	Branch if the result is Plus (>= 0)

Using these, we can test all sorts of things. One major use is to test for the end of a loop where an index, say .Y, has been set to a value. As the loop runs, .Y is decreased (DEY), then BEQ (when .Y = 0) to the next job, or BPL back around the loop if .Y is not zero. In just the same way, you could do one more loop by using BMI instead of BEQ. The variations are endless.

## LET'S WRITE SOMETHING

Let's put something on the screen where we can all see it.

```
LDY #$01 ;THIS IS SCREEN ASCII FOR 'A'  
STY $0400 ;THE FIRST SCREEN POSITION  
RTS ;BACK TO BASIC
```

That didn't hurt, did it. We put the screen ASCII for 'A' into .Y, then stored it in the top left of the screen. (on paper, at least)

## EXAMPLE 2

Surely we can do better, so lets put more than one letter up there.

```
LDY #$01 ;SCREEN 'A'  
STY $0400 ;ONTO THE SCREEN  
LDY #$02 ;SCREEN 'B'  
STY $0401 ;NEXT SCREEN POSITION  
LDY #$03 ;SCREEN 'C'  
STY $0402 ;NEXT SCREEN POSITION  
RTS ;BACK TO BASIC
```

Well, that works too, but it seems to be a long way to go to print, say, the alphabet. Are there better ways?

### EXAMPLE 3

There's (nearly) always a better way. Let's use all our cunning and make a loop to do it. Better still, we'll use indexed addressing to make it really swing.

```
LDY # $01 ; START A COUNTER
LDX # $09 ; AND ANOTHER HERE
HERE TYA   ; USE .Y AS CHAR
STA $0400,Y ; ONTO SCREEN
INY       ; NEXT LETTER AND POSITION INC Y
DEX       ; ONE LESS TO DO      DEC X
BPL HERE  ; IF +VE GO BACK TO 'HERE' NB DON'T NEED
RTS       ; IF NOT, ALL DONE    CMP TO BRANCH
```

The list introduces a new idea – did you notice 'HERE' on the third line? It's called a label, and roughly equates to a line number in BASIC. Labels are used for simplicity during program writing – they do not feature in the final ML. Whether it has a name (HERE) or is some address that the computer can find doesn't matter. It's simply a place in the program that can be identified.

What have we done? We have started the screen print with character 'A', and every time we go through the loop, the value stored in .Y is incremented by 1 (INY). A second counter is initialized with the value 9 in .X Each pass of the loop subtracts one from .X (DEX), and the value in .X is tested (BPL) to test whether to go back through the loop (branch to HERE) or to know the job is done and RTS. On each pass of the loop, the next character of the alphabet is printed to successive screen positions through the indexed counter based on \$0400 – 'A' in \$0401, 'B' in \$0402 and so on.

We now have a picture of the structure of a source file-

LABEL	OPC.OPERAND	COMMENT
	LDY #\$01	;START A COUNTER
	LDX #\$09	;AND ANOTHER
HERE	TYA	;USE .Y AS CHAR
	STA \$0400,Y	;ONTO SCREEN
etc		

We will return to this layout when we discuss program entry techniques later on in Topic 17 and Appendix A

# TOPIC 9 – Compares

Just as we can take branches depending on whether results are positive, negative, zero or not zero, so we can generate conditions to facilitate branching on the result of compares. Strictly speaking, BPL, BMI, BEQ and BNE are compares, but they are a special set of the family.

There are 3 compare opcodes – compare memory to .A, compare memory to .X and compare memory to .Y What sets this group aside is the method of operation. But first, the opcodes–

CMP	CoMPare to .A
CPX	ComPare to .X
CPY	ComPare to .Y

The mechanism of compares is not as you might expect. The compare operand is either a value, or the value contained in a memory location. The value is subtracted from .A, .X or .Y, but the value in the register is not changed. Nor is the result stored. The result of the compare can be negative, zero or positive (think about that for a minute), and CMP/CPX/CPY is usually followed by a branch, chosen depending on the desired effect of the compare.

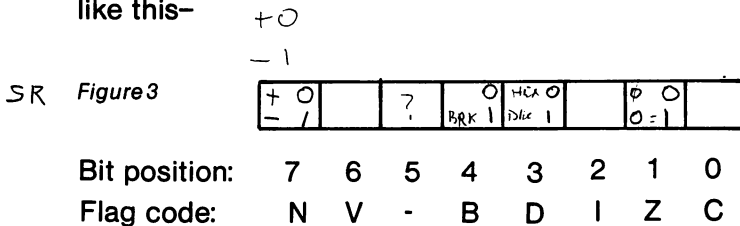


# TOPIC 10 — Flags

This topic has been held back until now, whereas in some texts it is introduced earlier. By now, many of the flag-setting instructions have been discussed without reference to flags. In simple programming, the flags are not so critical. But in more complex programming, or programs with more refinement, flag manipulation is very important. Pleasingly, it's really quite easy.

If you cast your mind back, we talked about flags in Topic 3, when we discussed the Status Register. There's nothing different about the Status Register – it's just a byte like any other, except that certain opcodes direct their enquiries there without any further prompting. So what does this byte do?

Just as we numbered bit positions in earlier discussions, the SR is indexed – with names this time. The byte (of the usual 8 bits) looks like this–



All looks mysterious, so some explanations.

**N – Bit 7 – The sign flag**, negative (N) if the bit is set (=1) or positive if 0.

**V – Bit 6 – The overflow flag**. It is set as a result of some mathematical operations. Little used.

Bit 5 is not assigned

*B – Bit 4 – The Break flag*, set if a break (BRK) is encountered.

*D – Bit 3 – Decimal mode flag*. When set, the computer uses BCD (binary coded decimal) mode, rather than hex mode for arithmetic operations.

*I – Bit 2 –* When set, the normal computer Interrupts (to check the keyboard etc) are defeated.

The next two registers are the most commonly used. Handily, they are at the end of the register where access is easiest.

*Z – Bit 1 – The Zero flag*. If the result of an operation is zero, Z is set. Otherwise, it is not set.

*C – Bit 0 – Carry flag*. Set when an operation produces a carry-over from one column of figures to the next more significant one, just as 6 + 5 in normal arithmetic produces a carry digit in the ‘tens’ column.

A number of opcodes work directly on these flags, to set or reset them. These are–

CLC	CLear the Carry flag (make it 0)
CLD	CLear the Decimal flag
CLI	CLear the Interrupt flag
CLV	CLear the oVerflow flag
SEC	SEt the Carry flag (make it 1)
SED	SEt the Decimal flag
SEI	SEt the Interrupt flag

Flags are important, and worthy of some study on their own. There is not time here to dissect their operation to any depth, but you need to be aware of their existence, and have an eye to their use.

To finish off, there are four Branch operations that we haven’t examined, as they needed some appreciation of flags before discussion. They are–

BCC	Branch if Carry (flag) Clear
BCS	Branch if Carry Set
BVC	Branch if oVerflow (flag) Clear
BVS	Branch if oVerflow Set

If the flag is clear ( $C=0$ ) when BCC is encountered, then the branch is taken. If set, the program merely passes straight through and on. Conversely, if BCS is the instruction and the flag is clear, the branch is not taken.

Similarly with BVC and BVS. The 'branch if clear' – 'branch if set' pattern from above is repeated, this time based on the overflow flag.



# TOPIC 11 — The Easy Ones

We need a short respite from complex thoughts, so let's look at 4 easily understood instructions.

The Stack(our LIFO pile of plates) can be used for many things, one of the more useful being to 'bank' the values in any or all of .A, .X or .Y before going off to do another task. Many programs in BASIC that have ML sections do this when the BASIC program calls (SYS XXXXX) the ML portion. How do we do this?

The most used instructions are—

PHA      PusH Accumulator value onto stack  
PLA      PuL Accumulator value off stack into .A

If all 3 registers are to be saved, there is a formula to follow that looks like—

```
PHA ;STORE .A
TYA ;PUT .Y IN .A
PHA ;STORE .Y
TXA ;PUT .X IN .A
PHA ;STORE .X
      ;REST OF PROGRAM
      ;
PLA ;GET ONE BACK
TAX ;PUT .A IN .X
PLA ;AND THE NEXT
TAY ;PUT .A IN .Y
PLA ;THAT'S ALL
```

Notice that we put the values onto the stack in order .A-.Y-.X and get them back in order .X-.Y-.A Provided you reverse the order on retrieval, this system is OK. Don't forget that .A must always be first in if you want to keep it.

NB

If you turn back to Topic 7, there was a mention of a better system for temporary storage of a useful parameter. Instead of using STX \$D7 to keep the variable, we could use PHA to put the value onto the stack, perform the INX, then PLA to get the value back.

The other instructions are to store the SR contents on the stack—i.e., the flags. The opcodes are—

- PHP      PusH Processor onto stack
- PLP      PuL Processor into SR

These need little explanation. But it is sometimes a way to preset the flags. If you set up .A, do a PHA then a PLP, the value you set up in .A finishes up in the SR.



# TOPIC 12 — Logical Operations

Many of you will not have come across Boolean functions before, so this will all be new to you. But don't worry – it is not difficult.

There are 4 instructions in this group–

AND	...AND!
BIT	This instruction tests memory bits
EOR	Exclusively OR memory with .A
ORA	Inclusively OR memory with .A

That sounds complicated, so a brief exposure to logical operations will help. Such operations do bit-by-bit comparisons of two values, locations, or whatever is specified. Bit 7 in one is compared to Bit 7 in the other, the result is stored, and so on to Bit 0. Bits are either 1 or 0, so a comparison could be comparing 1 with 1, 1 with 0, 0 with 1 or 0 with 0.

AND compares the 8 bits in .A with the specified value or location, bit by bit. If both bits in a comparison are 1 (1 AND 1) then the result is 1. Any other combination gives 0. The result is stored in .A An example–

83 is #\$53 – 01010011  
AND 112 is #\$70 – 01110000  
So (82 AND 112) is 01010000 which is 80 – \$50

The BIT instruction is an odd one to understand – two bits from memory (Bit 6 and Bit 7) are copied into the V and N flag. The Z flag is set if .A AND memory are not zero. Don't worry too much – this is for more advanced work.

EOR is also a bit-by-bit compare. This time, if either one or the other bit is 1, then the comparison result is 1. If both bits are 1 or both are 0, the result is 0. EOR is very handy for handling the Commodore

screen, particularly for reverse video. In Commodore code, the end bit (Bit 7) is off for normal display, and on for reverse. To reverse a character, Bit 7 has to be turned on-

```
LDA $0400 ;GET CHAR FROM SCREEN
EOR #$80 ;REVERSE IT
STA $0400 ;PUT IT BACK
```

'E' is 00000101  
Reverse it 10000000  
And get.. 10000101 which gives reverse 'E'

ORA is the last of the four opcodes to consider. During the comparison, if any of the bits being compared is 1, then the result is 1. An example is probably best-

```
Rev 'E' is 10000101
ORA #$80 10000000
And get.. 10000101 so it remains unchanged
```

A further example of AND and ORA can be used to blank the C64 screen during program load and save operations - this speeds up the process for disk, and is essential for cassette. Screen blanking is controlled through bit 4 of the VIC Control Register, location \$D011 (53265). When bit 4 is 'off', the screen is blanked. The routine looks like this-

```
LDA $D011 ;GET PRESENT VALUE
AND #$EF ;TURN BIT 4 OFF
STA $D011 ;PUT VALUE BACK
```

To turn the screen on again-

```
LDA $D011
ORA #$10 ;TURN BIT 4 ON
STA $D011
```

The effect of AND #\$EF is to leave all bits except bit 4 as they were. Bit 4 is made zero. In reverse, ORA #\$10 only affects bit 4. In BASIC, the routine would be-

```
POKE 53265, PEEK(53265) AND 239 = ON
POKE 53265, PEEK(53265) OR 16 = OFF
```

# TOPIC 13 — Addition and Subtraction

Arithmetic operations are not well implemented in Commodore computers. Commodore are not unique in this – it is common to all older microprocessors, but is improved in newer devices. But you will be able to get by (most of the time) without doing a vast amount of adding and subtracting. There are only two instructions in this group–

ADC      ADd with Carry  
SBC      SuBtract with Carry

With addition and subtraction, there can often be a carry – if the result exceeds 255 (#\$FF) or is less than 0 (#\$00), just as  $9 + 4$  produces a carry from units to tens in decimal addition. When this happens, the C (Carry) flag in the SR is set to 1. Of course, you may have already set C with a previous calculation, so it is usual to perform CLC (who remembers?) before using ADC. With subtraction, it is usual to do it the other way around– set C before an SBC.

The technique can be used to add large numbers, using C to indicate carry within each byte. Carry between bytes is handled by the 65XX for you. But let's stick with a simple addition–

```
CLC            ;JUST TO BE SURE
LDA $70       ;PUT MEM INTO .A
ADC #$16      ;ADD 22
STA $70       ;PUT IT BACK
SEC            ;SET THE FLAG
LDA $77       ;PUT MEM INTO .A
SBC #$04      ;TAKE AWAY 4
STA $77       ;PUT IT BACK
```

There are many and various ways of combining ADC and SBC with other instructions to produce quite powerful operations – and both ADC and SBC set and unset flags with gay abandon. Experiment !!!



# TOPIC 14 — Shifts and Things

Let's look at a byte – all 8 bits of it. Any value will do, so lets use decimal 42 – #2A

00101010

Lets move all the bits one place to the left–

01010100

This comes to #54, or 84 decimal. What do you notice?

Let's take #2A again, and shift the bits to the right this time–

00101010  
becomes 00010101

which is #15, or 21. The plot thickens, if you hadn't noticed. Without further ado, let's introduce the next set of instructions.

ASL	Arithmetic Shift Left
LSR	Logical Shift Right
ROL	ROtate Left
ROR	ROtate Right

The principal difference between these instructions is what happens to the bit that falls off the end, and what is put into the empty bit at the other 'end'.

ASL keeps the displaced bit in the Carry flag, and puts a zero into Bit 0.

C=1 ← 10010101  
00101010 ← 0

LSR is near enough to the reverse of ASL. Bit 0 goes into the Carry, and 0 is put into Bit 7

0 → 10010101 → C=1  
01001010

ROL is a nine-bit rotation. All bits are rotated one place left – the value in C goes into Bit 0, and then what was Bit 7 goes into the Carry flag.

C=1 ← 10010101 ← C=0  
00101010

ROR is another nine-bit rotation, the reverse of ROL. The Carry goes into Bit 7, and Bit 0 goes into the Carry flag.

### THE LAST ONES

*There are only three more opcodes. They are:*

SEI        which is SEt Interrupt  
CLI        which is CLear Interrupt  
RTI        which is ReTurn from Interrupt

*At this stage (and perhaps for a while yet) you won't need to worry about these. Just note their existence. When you are more experienced, you will learn to work with these useful commands.*

# TOPIC 15 — Other Addressing Modes

We have covered all the instructions, but there are some addressing modes that are a little more complex. I have left them until now, so that you will be in a better position to understand them.

Before starting in, there is a feature of the 65XX that you may already know. It concerns a peculiarity of 65XX architecture when storing numbers greater than # $\$FF$  (255 decimal) To store a number like 1036 ( $\$040C$ ) takes two bytes. One byte stores # $\$04$  and the other byte stores # $\$0C$ . But the 65XX likes to work back to front, so it stores # $\$0C$  in the lower byte and # $\$04$  in the higher byte. Almost without exception, you will see numbers stored this way in Commodore computers.

The main addressing mode that makes use of this feature is termed 'indirect', and it may be indexed or absolute. Let us consider indirect indexed by an example.

LDA ( $\$5E$ ),Y

translates as follows

LDA (the 2-byte address  $\$5F\$5E$ ),Y

so if  $\$5F$  contains # $\$04$  and  $\$5E$  contains # $\$03$  we have

LDA ( $\$0403$ ),Y

*LOAD A WITH CONTENTS OF 0403 + Y*

Note that the address is compounded from the two bytes (# $\$03$  and # $\$04$  in reverse order) THEN the value in .Y is added to determine the true location.

The value in .X can also be used, but the operation is different.

LDA ( $\$D7$ ),X

*Zero Page*

*5E*

*5E 03 260*

*5F 04 416*

*60*

The value in .X is added to the zero page address, to get a new address. This address, together with the next byte, is the location pointed to. For example, if \$D7 is #\$5C and .X is #\$02, then the instruction goes to the two bytes at \$D7 and \$D8. If \$D7 is #\$0A and \$D8 is #\$80, then the byte at \$800A is loaded into .A

eg. Zero Page.

D5 + LDA (D5, X)  
if X = 2.

D7 0A

D8 80

A is loaded with the contents of Address 800A.

# TOPIC 16 — To Finish Off

We have now looked at all the opcodes, and introduced the addressing concept. At this point, let's concentrate on reviewing what we know about addressing modes, and tabulate the individual modes.

The first mode was IMMEDIATE, where we loaded a variable into .A, .X or .Y.

```
LDA #$20 puts decimal 32 into .A
LDX #$45 puts decimal 69 into .X
```

The next mode was ABSOLUTE, where we loaded or stored information from or to an address. ABSOLUTE ZERO PAGE uses addresses between \$00 and \$FF – the first 256 bytes in the computer. ABSOLUTE (on its own) refers to any other address that the computer can access.

```
LDA $B000
STA $A7
```

After that came INDEXED mode – ZERO PAGE INDEXED by .Y or .X (zero page meaning the same as above) and ABSOLUTE INDEXED, for any other address.

```
LDA $B147,Y
STA $0400,X
LDA $D7,Y
```

Branches followed – a mode on its own, termed RELATIVE because the jump 'distance' is related to the bytes being executed.

IMPLICIT was straightforward – INY adds one to .Y, DEX takes 1 from .X, PHA pushes the accumulator onto the stack, and so on. The opcodes stand alone, implying their action without need for any more explanation.

The INDEXED mode was next. In this mode, the computer calculates its action using a 65XX shorthand convention:-

LDA (\$42),Y

gets the bytes from \$42 and \$43, applies the 65XX convention of low byte-high byte and loads A with the contents of that address- in this case, offset by .Y In this case, the mode is termed INDIRECT INDEXED.

As a shorthand guide, the following table sets out the major modes:-

IMMEDIATE	LDA #\$00
ABSOLUTE	LDA \$C000
ABSOLUTE ZERO PAGE	LDA \$2A
ABSOLUTE INDEXED	LDA \$B000,Y
ZERO PAGE INDEXED	LDA \$20,Y
RELATIVE	BNE \$1024
IMPLIED	INY
INDIRECT INDEXED	LDA (\$2A),Y
ACCUMULATOR	ASL A

There is one other mode, not covered in the main text. INDIRECT (sometimes termed ABSOLUTE INDIRECT) means that the program goes straight to an address, finds the contents of the address and the byte following it, and goes to that place to continue.

Some texts expand these to indicate that some modes work only with certain parameters, one example being to indicate whether .X, .Y or both can be used as indices. However, the principal modes are listed above.

These names are used somewhat flexibly by different writers. This can cause some confusion and user difficulty, but if you learn to say the mode to yourself each time you use the various instructions, they will form a distinct pattern in your mind and ease the programming task.

# TOPIC 17 — Editors and Assemblers

Programming in ML is nothing like BASIC programming. With BASIC, you are dealing in quasi-English, working with a dynamic screen, and you can insert, delete, move, renumber, rename and otherwise fiddle to your heart's content – all nicely protected by a BASIC interpreter that does all sorts of convenient things for you, even down to telling you where your program ground to a halt. Not so ML – there is no tidy entry routine directly available for use. You have to look for help.

Up to now, we have been dealing in labels, opcodes and operands, which are reasonably easy to read by us, but which are of no direct value to the computer.

At the outset, it was suggested that you get a copy of an extended monitor – SUPERMON or some such. These programs contain a 'tiny assembler', which translates opcodes and operands into the bytes that are understood by the computer.

No, the computer doesn't understand things like LDA #\$01 – but it does understand a translation of it. Who or what does the translation?

To understand that, let us go through the process of assembly. For small programs, the tiny assemblers are adequate, but you should have a rough draft of your program down on paper. So, tiny assemblers are useful, but restricted. Fire up your monitor program, and let's practice. Say it's SUPERMON, for the example.

SUPERMON (like most monitors and extensions) prompts you for an entry with a '.' To assemble with SUPERMON, type 'A' after the '.', a space, then the start address (in hex) of the first instruction, then the opcode and operand.

.A 1000 LDA #\$01 <Carriage Return>

SUPERMON will promptly translate this into hex bytes, store them and print them thus

```
.A 1000 A9 01
```

and prompt you for the next instruction by printing the next address (\$1002 in this case).

```
.A 1000 A9 01
```

```
.A 1002
```

You continue entering the program, and simply press RETURN (<CR>) when entry is finished.

For larger jobs, some form of Assembler package is needed. There are a number of these for disk users, the best known being MAE (disk only) or ASSM/TED (for disk and cassette users) and the CBM Assembler packages for the CBM/PET series and the C64. These will be discussed later. Each package contains programs to edit, assemble and load ML programs. Let's look at each of these functions.

An editor allows the user to enter instructions in simple form, something like entries in SUPERMON as described above. Usually the entry looks something like a BASIC program – line number and so on–

```
10 LDA #$01  
20 STA $22 ;COUNTER FOR PAGES  
30 LDA #$42  
40 STA $23 ;66 LINES PER PAGE
```

Editors usually permit comments (“REMs”) as well (the ‘;’ is translated to REM), so that the function of each line can be noted. Files, saved either on disk or cassette, contain sufficient instructions to create the desired program, and are usually termed ‘source files’, the name indicating that the file contains the original material for the eventual program.

The next task is to assemble (translate, interpret, convert) the source file. The assembler converts all opcode/operand combinations

to hex, in a form understood by the computer. During this phase, the assembler checks for syntax validity, but not whether the program will work or fail. That's up to you!! Also during assembly, you may create another file, containing the hex bytes only. This file is usually termed the 'object code', 'code' as it is all in hex.

Finally, the object code is loaded, usually using another routine that reads the object code and stores it in memory. At that point, the program is ready for the final test – run time.

Sometimes, the ML bytes that have been loaded into the computer are 'saved' as a file of their own – termed a 'binary file'. This binary file can later be loaded directly back into the computer without need for an intermediate assembler program. At run time, this can be very convenient.

We will use all of these features to provide some knowledge of their use. Most assembler packages are documented (for better or worse – usually ambiguously) so you can get further guidance there. If all else fails, you could always contact a more experienced user.

## **ASSEMBLER SUITES**

As noted above, the principal packages are MAE (Macro Assembler Editor) and the CBM Assembler Packages. The former is a commercial product, with versions for CBM/PET and C64. The package is largely self-contained, very powerful, with some very nice features that make its use a pleasure. In addition, there is a group in the USA (ATUG – ASSM/TED Users Group) who promote/adapt/improve MAE, with membership open to MAE owners. If MAE has a drawback, it is that source files are restricted to about 1K each. As it is possible to 'chain' (connect) source files to form very large program files, this is not a total limitation.

The CBM packages are Public Domain Software, available from Commodore dealers at a nominal price (\$60 when last checked) or your local User Group for a copying fee. The package contains 4 discrete program groups for disk operation - Editor, Assembler, Loader and Cross Reference, each of which has to be individually loaded as the need arises. The software lacks the sophistication of MAE, but overall is very useable. Documentation is adequate but not over-helpful.

Because of lower cost, the CBM Assembler Packages (3032/4032/8032, and C64 versions) are attractive, so it may be the starting point for would-be ML programmers. Appendix A sets out an outline technique for using the package.

# AND IN CONCLUSION

By now, you should have a general outline of ML programming. The book was not intended to provide an in-depth, exhaustive volume designed to lead straight to a degree in computer science. If I have not confused you too thoroughly, and you still want to do some self-motivated study, then you are on your way to ML proficiency.

A second volume is planned, to bring together some useful ideas and routines that will form a basis for your ML program library.

*Paul Blair*



# APPENDIX A

## USING THE COMMODORE 64 PACKAGE

The Commodore package is designed for use with a disk drive. There is no cassette version.

The package consists of 4 main elements –

EDITOR  
ASSEMBLER  
LOADER  
CROSS REFERENCE

### **EDITOR:-**

This is a ML program that loads into memory (LOAD "EDIT.C64",8,1) at \$C000 (49152), and is activated by a SYS to that location. The EDITOR provides for preparation of input files, and helps by some additional features. These include –

AUTO – to provide automatic line number entry.  
CHANGE/FIND – to find and/or amend any characters.  
DELETE – erase any range of lines.  
FORMAT – tab source file to the screen.  
GET – the EDITOR's 'Load'.  
KILL – turn off the EDITOR.  
NUMBER – renumber the file.  
PUT – the EDITOR's normal 'Save' to disk.  
CPUT – the EDITOR's compressed 'Save'.

You will recall from Topic 8 that the format used for file entry is –

LINE NUMBER (LABEL) OPCODE OPERAND (COMMENT)

Titles in parentheses are optional. The operand is not required when using implied opcodes (TAX INY etc.)

So, let's start with a program. Assume that you have drafted out your program on paper, then loaded and SYS'ed the EDITOR into life. What now?

We start typing –

```
100 START LDX #$00           ;ZERO INTO COUNTERS
110 STX $B2                 ;COUNTER 1
120 STX $B3                 ;COUNTER 2
130 TXA                    ;MOVE TO .A
140 TAY                    ;SET INDEX TO 0
150 TITLE LDA (HEADER), Y   ;GET A CHARACTER
160 AND #$7F                ;CORRECT FOR SCREEN
170 JSR $FFD2               ;ROM PRINT ROUTINE
180 INY                    ;ADD 1 TO INDEX
190 CPY #$20                ;DONE 32 YET?
200 BNE TITLE               ;LOOP BACK
210 OPEN .....
```

The program is entered somewhat like a BASIC program – line number and text, except that the EDITOR program is handling things now, not the BASIC interpreter in the computer. In fact, while in EDITOR mode, it's not a good idea to try any BASIC programs. KILL the EDITOR and go back to BASIC if you want to do that.

The physical layout shown here is deliberate. The fields are separated by a space, essential for the ASSEMBLER program. The exception is where a label is not used, and you should leave TWO spaces after the line number. This is so that the FORMAT command will work. The ASSEMBLER can cope with this variation, so the extra space causes no problems.

When you have typed in the program, the next step is to store it on disk. PUT is the command, so we –

PUT“0:FILENAME”

There are options here. The drive number can be omitted, and the EDITOR will default to the last drive used (not a problem with single disk drives!) and store your text (without line numbers) as a SEQ file. I have used the ‘save with replace’ (@0:) without any problems on DOS 1 and 2 – but you should follow your own practice here, as there have been assertions as to the unreliability of this syntax. PUT also permits line number ranges to be selected if desired, as you use LIST in BASIC, to file your text selectively.

A variation of PUT is CPUT. The effect is the same as PUT, except that unnecessary spaces are not written to file. The idea of this is to speed up the Assembler process.

The opposite of PUT is GET – not BASIC GET but EDITOR GET. Syntax is easy

GET “FILENAME”

without drive number will do. The sequential file will be loaded, and assigned line numbers starting from 1000 and going up in tens. If need be, you can now edit your program, then PUT it back on disk.

### **ASSEMBLER :-**

More ML, this time residing at the usual BASIC start. (LOAD“ASM.C64”,8) The ASSEMBLER can co-reside with the EDITOR, which is safely at \$C000. In fact, to return to the EDITOR after using ASSEMBLER, simply do SYS 49152.

To activate the ASSEMBLER, type RUN and you’re in business. You will be asked some questions – note that no quotation marks need be typed for entries in this phase. Do you want an object file, ready to be loaded into the computer? If so, on which drive and with what name? If no file is needed (say during development) press <CR>. The next question concerns hardcopy. If required then <CR> again. If not, type “N” with <CR>. If you said “N” to an object file, you will be asked if you want cross reference files. Respond as your needs dictate. Finally, you need to enter the name of the source file to be assembled. The drive number can be specified, but is not

mandatory. If all has been correctly entered, then ASSEMBLER goes to it.

ASSEMBLER makes two passes, looking for labels, jumps and branches on the first pass, then computing all values on the second pass. At the end (hopefully) you will be told "ERRORS = 0000" or (less hopefully) the number that were found. Remember, these are syntactical errors – the 'nil errors' message does not guarantee the logic of your program!

Assuming that you requested an object file, the next step is to load the information into the computer. Before doing so, it is advisable to completely reset the computer.

At this point, there is one step that you will probably want to take. If you are going to load the object file into your computer, it is likely that you will want to save the ML bytes for later. To do this, you will need a monitor program to provide the necessary save function. Select a monitor unlikely to be affected by the loading process (a monitor in \$C000 is usually pretty safe) and get it in place. Now, to continue....

Depending on where you intend to store the program when loaded, you have a choice of LOADERS, LO-LOAD and HI-LOAD. Check this table for usage –

PROGRAM TO BE LOADED TO	USE THE x-LOADER
LOW MEMORY	HIGH
MID MEMORY	LOW or HIGH
HIGH MEMORY	LOW

LOADER	TO RUN
LOW	RUN
HIGH	SYS 51200

The loaders will firstly ask for an offset. If your program is written to reside at \$1000 but you want it temporarily stored at \$3000, then the offset is \$2000. You might not see the point of offsets just yet, but the time will arrive.....

For now, ignore the question by a carriage return (<CR>) and go on to the real task – loading the object file. Enter the filename, <CR>, and all being well you will see the starting address (in hex) of the actual loading place on the screen. This will be followed by a series of dots, indicating the progress of the load. Finally, the end address of the actual storage location is printed on the screen. Remember or write down these addresses for the next step.

The message 'END OF LOAD' signifies the end of the load phase. Type 'NEW' to reset pointers, and then (before doing anything else) SAVE the program. Jump into the monitor at the address you stored it, then

```
. S'd:filename',08,aaaa,bbbb <CR>
```

where 'd' is the disk drive number, 'aaaa' is the start of your program and 'bbbb' is ONE BYTE PAST the end of it. If you don't go one byte past, you will lose the last byte of your opus.

That's it. You now have a copy in the computer, and a copy on disk. What you do next is up to you.

If you have any problems with this routine, have a look at this diagram, which is a flowchart of the main points written so far. If this all looks a bit long winded, then consider what life was like before assembler packages were developed and all programs were hand written, assembled and entered!!

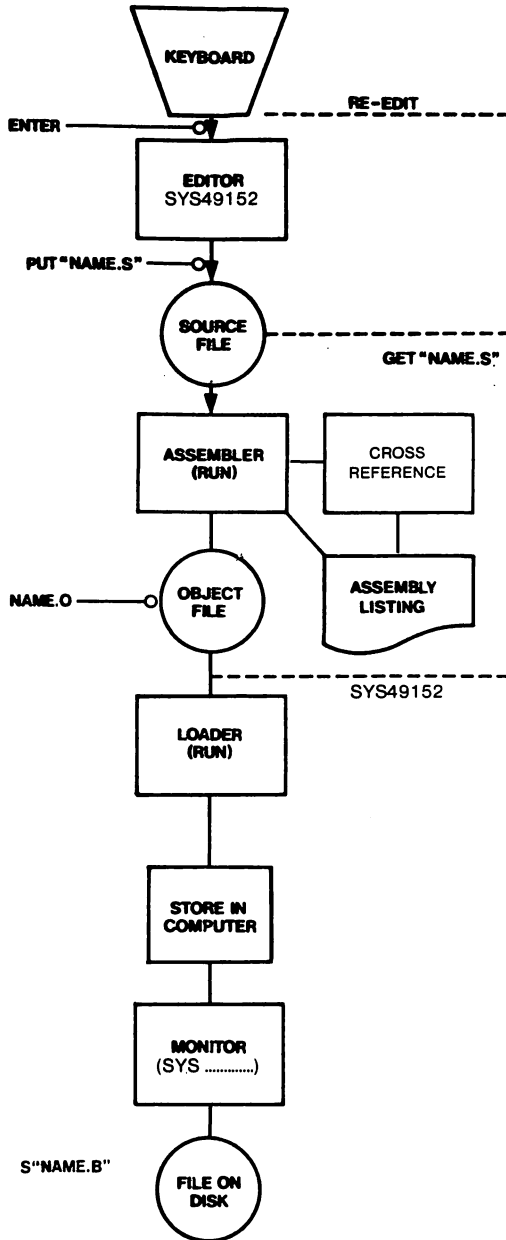
A useful feature now implemented in all Commodore Assembler Packages is the cross reference option. At the start of assembly, you may choose to create an object file, or a series of cross reference files. You can have one or the other, but not both.

If you choose the cross reference option, the Assembler will create a series of files on disk during run time. When the Assembler has finished, load CROSSREF64 and type RUN <CR>. You may elect

to direct output to the screen or printer. The result is a listing that shows all labels used in the source files, and the Assembler output line numbers where the labels occur. This is particularly useful for later reference, trouble shooting or unravelling programs that are more complex.

That completes a run through the principal technique required to use the Commodore Assembler package. The documentation explains a wider range of sub-options that you can use in the development process.

# ASSEMBLY PROCEDURE FLOWCHART



# NAMING CONVENTIONS

We have noted that the assembly process creates three files (they are usually called 'files' for some obscure reason) the original source file, the object file, and the ML file. To ease the task of subsequent use by you or others, it is convenient (but not mandatory) to use some convention to identify which is which. One convention that has become widespread is to use the following suffixes to file names:-

- .SRC or .S for source files
- .OBJ or .O for object files
- .BIN or .B for ML (binary) files

To some, the SOB initials sum up ML programming very succinctly!

ML is sometimes converted into a BASIC loader program. Each byte of ML is kept in DATA statements, which the BASIC program READs and the values are POKE'd into the appropriate memory location. If you do this, the BASIC program could be given the suffix .BAS- but be careful not to use .B for binary AND Basic files.

# APPENDIX B

## A SAMPLE GRAPHICS PROGRAM

Appendix B is presented for study in conjunction with the introduction to programming presented in this book.

As the printout indicates, the program has been written for a Commodore 64 computer.

As the program notes, it is long on detail and (somewhat) short on sophistication. This is deliberate, because it will permit a better grasp of the fundamental steps used in the program. To tease you, most of the routines could be accomplished in fewer bytes – but I'll leave the investigation and solution of that to you. If you don't want to type in the program, disk and cassette copies are available from where you purchased this book or through KIM BOOKS.

Read on. The more you read/study other peoples programs, the faster your abilities in ML will grow.

# BASIC LOADER

If you would like to load the program without having to go through the assembly process, here is a BASIC loader that will do the job for you. I have given the listing here in short form, without line numbers and any pretty screen routines. You may add them to taste.

BASIC LOADER FOR C64 ML PAUL BLAIR – COMPUCHART

PRINT'LOADING C64ML PROGRAM'

PRINT'PLEASE WAIT'

S=49152 : F=50263 : FOR I = S TO F

READ A : POKE I,A : NEXT

PRINT'LOADING COMPLETE'

PRINT'AND ROUTINES INITIALIZED'

PRINT

PRINT'(C) KIMBOOKS 1984' : SYS 49152 : NEW

DATA 169,76,133,115,169,26,133,116,169,192,133,117,96,  
169,230,133

DATA 115,169,122,133,116,169,208,133,117,96,142,255,207,  
186,189,1

DATA 1,201,140,208,16,189,2,1,201,164,208,20,230,122,208,2

DATA 230,123,76,76,192,201,230,208,7,189,2,1,201,167  
,240,236

DATA 174,255,207,230,122,208,2,230,123,76,121,0,160,0,  
177,122

DATA 201,64,240,3,76,121,0,230,122,208,2,230,123,177,  
122,230

DATA 122,208,2,230,123,201,66,208,3,76,223,192,201,67,  
208,3

DATA 76,63,193,201,68,208,3,76,110,193,201,69,208,3,76,175

DATA 193,201,70,208,3,76,240,193,201,73,208,3,76,29,  
194,201

DATA 74,208,3,76,65,194,201,75,208,3,76,112,194,201,76,208

DATA 3,76,129,194,201,77,208,3,76,186,194,201,78,208,3,76

DATA 225,194,201,82,208,3,76,251,194,201,83,208,3,76,59,  
195  
DATA 201,84,208,3,76,116,195,201,85,208,3,76,137,195,201,  
87  
DATA 208,3,76,188,195,201,90,240,3,76,8,175,76,13,192,169  
DATA 6,133,76,32,155,183,134,77,169,247,229,77,133,75,  
32,241  
DATA 183,134,140,32,241,183,142,134,2,165,140,201,8,48,  
26,233  
DATA 8,133,140,169,160,145,75,32,45,193,24,165,75,233,  
39,133  
DATA 75,165,76,233,0,133,76,208,224,170,189,37,193,145,  
75,32  
DATA 45,193,76,115,0,100,111,121,98,248,247,227,160,169,  
220,69  
DATA 76,133,76,173,134,2,145,75,169,220,69,76,133,76,96,32  
DATA 155,183,169,4,133,76,169,39,134,77,229,77,105,79,  
133,75  
DATA 169,32,145,75,24,165,75,105,40,133,75,165,76,105,  
0,133  
DATA 76,201,7,208,235,165,75,201,32,48,229,76,115,0,160,40  
DATA 132,142,160,0,132,140,160,191,162,7,134,141,134,143,  
177,140  
DATA 145,142,32,162,193,177,140,145,142,32,162,193,136,  
192,255,208  
DATA 237,202,224,3,208,228,160,39,169,32,145,140,136,16,  
251,76  
DATA 121,0,169,220,69,143,133,143,169,220,69,141,133,141,  
96,169  
DATA 2,162,8,168,32,186,255,32,13,196,152,162,0,160,207,32  
DATA 189,255,32,192,255,162,2,32,198,255,32,207,255,32,  
210,255  
DATA 36,144,112,17,201,13,208,242,173,141,2,41,1,208,  
249,165  
DATA 197,201,63,208,229,32,204,255,169,2,32,195,255,76,  
115,0  
DATA 32,155,183,134,75,32,241,183,142,134,2,160,0,132,  
140,162  
DATA 4,165,75,134,141,145,140,32,168,193,173,134,2,145,  
140,32  
DATA 168,193,165,75,200,208,238,232,224,8,208,231,96,162,  
4,134

DATA 76,160,0,132,75,32,155,183,142,134,2,166,76,177,75,73  
DATA 128,145,75,32,45,193,200,208,244,232,134,76,224,8,  
208,237  
DATA 96,32,13,196,169,40,132,77,229,77,133,77,169,32,  
32,210  
DATA 255,198,77,198,77,165,77,201,1,16,241,160,0,185,0,207  
DATA 240,6,32,210,255,200,208,245,169,13,32,210,255,76,  
115,0  
DATA 169,49,141,20,3,169,234,141,21,3,169,16,133,252,  
76,121  
DATA 0,162,4,160,0,132,140,132,142,169,195,133,143,134,  
141,230  
DATA 143,230,143,169,0,133,142,177,142,145,140,32,168,  
193,230,143  
DATA 177,142,145,140,32,168,193,198,143,200,208,235,232,  
224,8,208  
DATA 220,173,254,207,141,33,208,76,121,0,32,155,183,134,  
76,32  
DATA 241,183,134,75,160,0,132,140,162,4,165,75,134,141,  
177,140  
DATA 197,76,208,4,165,75,145,140,200,208,243,232,224,8,  
208,236  
DATA 96,32,158,173,32,247,183,72,169,36,32,210,255,104,  
32,58  
DATA 196,152,32,58,196,169,13,32,210,255,96,120,169,14,  
141,20  
DATA 3,169,195,141,21,3,169,1,133,78,88,76,122,194,165,197  
DATA 197,254,240,9,133,254,169,16,133,253,76,49,234,201,  
255,240  
DATA 249,165,253,240,4,198,253,208,241,198,78,208,237,  
169,4,133  
DATA 78,169,64,133,197,169,2,133,205,208,223,173,33,208,  
141,254  
DATA 207,162,4,160,0,132,140,132,142,169,195,133,143,134,  
141,230  
DATA 143,230,143,169,0,133,142,177,140,145,142,32,168,  
193,230,143  
DATA 177,140,145,142,32,168,193,198,143,200,208,235,232,  
224,8,208  
DATA 220,76,121,0,32,155,183,134,75,32,241,183,134,76,  
166,75

DATA 164,76,24,32,240,255,76,115,0,162,0,134,142,162,  
40,134  
DATA 140,162,4,134,143,134,141,160,0,177,140,145,142,32,  
162,193  
DATA 177,140,145,142,32,162,193,200,208,239,232,224,8,  
208,228,169  
DATA 32,162,40,157,199,7,202,208,250,76,121,0,173,33,  
208,141  
DATA 253,207,162,4,160,0,132,140,132,142,169,195,133,  
143,134,141  
DATA 230,143,230,143,177,142,133,75,177,140,145,142,  
165,75,145,140  
DATA 32,168,193,230,143,177,142,133,75,177,140,145,142,  
165,75,145  
DATA 140,198,143,32,168,193,200,208,219,232,224,8,208,  
208,173,254  
DATA 207,141,33,208,173,253,207,141,254,207,76,121,0,32,  
115,0  
DATA 201,34,240,3,76,8,175,160,1,177,122,240,247,201,  
34,240  
DATA 6,153,255,206,200,208,242,24,169,0,153,255,206,152,  
101,122  
DATA 133,122,169,0,101,123,133,123,136,96,72,74,74,74,  
74,201  
DATA 10,144,2,105,6,105,48,32,210,255,104,41,15,201,10,144  
DATA 2,105,6,105,48,76,210,255

C64ML.SRC

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00001	0000		*****			*****
00002	0000		*			*
00003	0000		*		'ML ON THE C64'	*
00004	0000		*			*
00005	0000		*		ML PROGRAMMING EXAMPLES	*
00006	0000		*			*
00007	0000		*		COMMODORE 64	*
00008	0000		*		40 COLUMN SCREEN	*
00009	0000		*			*
00010	0000		*		ADAPTED FOR CBM/PET	*
00011	0000		*		-PAUL BLAIR	*
00012	0000		*			*
00013	0000		*		ADAPTED FOR C64	*
00014	0000		*		-PAUL BLAIR	*
00015	0000		*		-PETER FLETCHER	*
00016	0000		*			*
00017	0000		*****			*****
00019	0000					THIS PROGRAM PROVIDES SOME USEFUL
00020	0000					EXAMPLES OF SHORT ML ROUTINES,
00021	0000					MAINLY FOR SCREEN HANDLING.

```

00022 0000 ;THE ROUTINES ARE AVAILABLE IN
00023 0000 ;BOTH DIRECT AND PROGRAM MODE.
00024 0000 ;
00025 0000 ;EACH ROUTINE IS COMMENTED TO
00026 0000 ;DESCRIBE ITS METHOD OF OPERATION.
00027 0000 ;
00028 0000 ;THE ROUTINES ARE NOT SOPHISTICATED,
00029 0000 ;THEY HAVE BEEN WRITTEN TO BE
00030 0000 ;ILLUSTRATIVE. THERE ARE MORE
00031 0000 ;COMPLEX AND INTRICATE WAYS OF
00032 0000 ;DOING THEM - YOU SHOULD BE ABLE
00033 0000 ;TO IDENTIFY WAYS OF MAKING
00034 0000 ;IMPROVEMENTS.

```

```

00036 0000 ;THE ROUTINES USE THE '@' KEY TO
00037 0000 ;PREFACE THE COMMANDS - EG @R WILL
00038 0000 ;TURN ON REPEAT FOR ALL KEYS, @K
00039 0000 ;WILL TURN IT OFF. @E "XXXXXX"
00040 0000 ;READS SEQUENTIAL FILE "XXXXXX"
00041 0000 ;DIRECT TO THE SCREEN WITHOUT
00042 0000 ;DISTURBING ANY PROGRAM IN MEMORY.
00043 0000 ;
00044 0000 ;ALL COMMANDS ARE NOTED IN THE PROGRAM
00045 0000 ;LISTING, TOGETHER WITH THE CORRECT
00046 0000 ;SYNTAX FOR USE. HAVE FUN!!

```

ADDRESSES

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00048	0000		*****			
00049	0000		* DEFINE MACHINE LABELS	*		
00050	0000		*****			
00051	0000					
00052	0000		;---ZERO PAGE---			
00053	0000					
00054	0000		USRO	=\$4B		; THREE LOCATIONS
00055	0000		USR1	=\$4C		;.. FOR TEMPORARY
00056	0000		USR2	=\$4D		;.. STORAGE
00057	0000		REPDL	=\$4E		
00058	0000		INDEX1	=\$8C		; FOUR LOCATIONS
00059	0000		INDEX2	=\$8D		;.. TO ACT
00060	0000		INDEX3	=\$8E		;.. AS POINTERS
00061	0000		INDEX4	=\$8F		;.. TO PLACES
00062	0000		CHRGET	=\$73		;GET NEXT CHAR FROM BASIC
00063	0000		CHRG7	=CHRGET+7		
00064	0000		CHRG8	=CHRGET+8		
00065	0000		CHRGOT	=\$0079		; REGET CHAR FROM BASIC
00066	0000		SATUS	=\$90		; VARIABLE ST
00067	0000		LSTX	=\$C5		; WHICH KEY HAS BEEN PUSHED?
00068	0000		BLNCT	=\$CD		; BLINK COUNTER FOR CURSOR
00069	0000		TEMP1	=\$FC		; 3 STORAGE SPOTS
00070	0000		TEMP2	=\$FD		
00071	0000		TEMP3	=\$FE		
00072	0000					
00073	0000		;---OTHER PAGES---			

```

00074 0000 ;
00075 0000 ; CHARACTER COLOUR
00076 0000 ; HAS SHIFT KEY BEEN PUSHED?
00077 0000 ; IRQ INTERRUPT VECTOR LOW
00078 0000 ; IRQ INTERRUPT VECTOR HIGH
00079 0000 ;
00080 0000 ; SCREEN = $0400
00081 0000 ;
00082 0000 ; ---C64 ROM 2.0 ROUTINES---
00083 0000 ;
00084 0000 ; READY = $A474
00085 0000 ; FRMNUM = $AD9E
00086 0000 ; SNERR = $AF08
00087 0000 ; GTBYTE = $B79B
00088 0000 ; COMBYT = $B7F1
00089 0000 ; GETADR = $B7F7
00090 0000 ; COLTMP = $CFFD
00091 0000 ; POP = $CFFF
00092 0000 ; SCCOL = $D021
00093 0000 ; KEY = $EA31
00094 0000 ; CLOSE = $FFC3
00095 0000 ; SETLFS = $FFBA
00096 0000 ; SETNAM = $FFBD
00097 0000 ; OPEN = $FFC0
00098 0000 ; CHKIN = $FFC6
00099 0000 ; CLRCHN = $FFCC
00100 0000 ; INCHR = $FFCF
00101 0000 ; OUTCH = $FFD2
00102 0000 ; PLOT = $FFF0
; COLOUR = $0286
; SFST = $028D
; CINV1 = $0314
; CINV2 = $0315
; SCREEN = $0400
; ---C64 ROM 2.0 ROUTINES---
;
; READY = $A474
; FRMNUM = $AD9E
; SNERR = $AF08
; GTBYTE = $B79B
; COMBYT = $B7F1
; GETADR = $B7F7
; COLTMP = $CFFD
; POP = $CFFF
; SCCOL = $D021
; KEY = $EA31
; CLOSE = $FFC3
; SETLFS = $FFBA
; SETNAM = $FFBD
; OPEN = $FFC0
; CHKIN = $FFC6
; CLRCHN = $FFCC
; INCHR = $FFCF
; OUTCH = $FFD2
; PLOT = $FFF0
; CHARACTER COLOUR
; HAS SHIFT KEY BEEN PUSHED?
; IRQ INTERRUPT VECTOR LOW
; IRQ INTERRUPT VECTOR HIGH
; C64 SCREEN RAM
; BACK TO BASIC
; EVALUATE EXPRESSION
; SYNTAX ERROR
; CHRGET + GET BYTE
; EVALUATE FORMULA
; FAC TO INTEGER
; TEMP SCREEN COL
; KEY POP
; SCREEN COLOUR
; INTERRUPT
; CLOSE FILE IN '.A'
; SET LA,FA,SA
; SET FILENAME PTR
; OPEN LOGICAL FILE
; JUMP TO SET INPUT DEVICE
; CLOSE CHANNELS
; INPUT A CHAR
; OUTPUT A CHAR
; PLOT CURSOR POS'N

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00104	0000		;		-----	
00105	0000		*	=\$C000	;	START AT SYS49152
00106	0000		;		-----	
00107	0000		;			
00108	0000		;			
00109	0000		*	ENABLE ROUTINE :	SYS 49152	*
00110	0000		;			
00111	0000		;			
00112	0000	A9 4C	LDA	#\$4C	;	POINT TO MAIN ENTRY
00113	0002	85 73	STA	CHRGET		
00114	0004	A9 1A	LDA	#<ENTRY		GET LOW ADDRESS
00115	0006	85 74	STA	CHRGET+1		STORE IT
00116	0008	A9 C0	LDA	#>ENTRY		GET HIGH ADDRESS
00117	000A	85 75	STA	CHRGET+2		STORE IT
00118	000C	60	RTS			BACK TO BASIC
00119	000D		;			
00120	000D		;			
00121	000D		*	DISABLE ROUTINE :	SYS 49165	*
00122	000D		;			
00123	000D		;			
00124	000D		;			
00125	000D		;			
00126	000D	A9 E6	LDA	#\$E6		
00127	000F	85 73	STA	CHRGET		
00128	0011	A9 7A	LDA	#CHRG7		
00129	0013	85 74	STA	CHRGET+1		

```

00130 C015 A9 D0 LDA #$D0
00131 C017 85 75 STA CHRGET+2
00132 C019 60 RTS
00133 C01A
00134 C01A ;*****
00135 C01A ;* MAIN ENTRY *
00136 C01A ;*****
00137 C01A ;
00138 C01A ENTRY STX POP
00139 C01D TSX
00140 C01E LDA $0101,X
00141 C021 CMP #$8C
00142 C023 BNE MAYBE
00143 C025 LDA $0102,X
00144 C028 CMP #$A4
00145 C02A BNE NOT
00146 C02C E6 7A INC CHR7
00147 C02E D0 02 BNE NEXT1
00148 C030 E6 7B INC CHR8
00149 C032 ;
00150 C032 4C 4C C0 ; NEXT1 JMP KEY01
00151 C035 ;
00152 C035 MAYBE
00153 C037 D0 07 CMP #$E6
00154 C039 BD 02 01 BNE NOT
00155 C03C C9 A7 LDA $0102,X
00156 C03E F0 EC CMP #$A7
00157 C040 AE FF CF BEQ OK
00158 C043 E6 7A INC CHR7
00159 C045 D0 02 BNE LAST1
; GET KEY BACK
; ADD 1 TO POINTER LOW
; STORE IT
; GET STACK POINTER
; GET CHAR
; DIRECT MODE?
; MAYBE
; GET NEXT CHAR
; PRG MODE
; NO
; ADD 1 TO POINTER LOW
; ADD 1 TO POINTER HIGH
; OK TO GO

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00160	C047	E6 7B		INC	CHRG8	; ADD 1 TO POINTER HIGH
00161	C049					
00162	C049	4C 79 00	LAST1	JMP	CHRGOT	; REGET KEY
00163	C04C					
00164	C04C					
00165	C04C					
00166	C04C					
00167	C04C					
00168	C04C	A0 00	KEY01	LDY	#\$00	; SET INDEX TO 0
00169	C04E	B1 7A		LDA	(CHRG7), Y	; GET CHAR
00170	C050	C9 40		CMP	#\$40	; START WITH '@'?
00171	C052	F0 03		BEQ	KEY02	; YES, CONTINUE
00172	C054	4C 79 00		JMP	CHRGOT	; NO, REGET CHAR
00173	C057					
00174	C057	E6 7A	KEY02	INC	CHRG7	; NEXT PLEASE
00175	C059	D0 02		BNE	KEY03	; CHRG7 < 0
00176	C05B	E6 7B		INC	CHRG8	; CARRY NEEDED
00177	C05D					
00178	C05D	B1 7A	KEY03	LDA	(CHRG7), Y	; GET NEXT CHAR
00179	C05F	E6 7A		INC	CHRG7	; COUNT ON
00180	C061	D0 02		BNE	KEY04	; CHRG7 STILL < 0
00181	C063	E6 7B		INC	CHRG8	; SO INCREMENT CHRG8
00182	C065					
00183	C065	C9 42	KEY04	CMP	#\$42	; 'B'?
00184	C067	D0 03		BNE	KEY05	; NO, KEEP SEARCHING
00185	C069	4C DF 00		JMP	BO1	; GO TO 'B' ROUTINE

```

00186 C06C          C9 43          ; KEY05      CMP #$43
00187 C06C          D0 03          BNE KEY06
00188 C06E          4C 3F C1      JMP C01
00189 C070          C9 44          ; KEY06      CMP #$44
00190 C073          D0 03          BNE KEY07
00191 C075          4C 6E C1      JMP D01
00192 C077          C9 45          ; KEY07      CMP #$45
00193 C07A          D0 03          BNE KEY08
00194 C07C          4C AF C1      JMP E01
00195 C07E          C9 46          ; KEY08      CMP #$46
00196 C081          D0 03          BNE KEY09
00197 C083          4C F0 C1      JMP F01
00198 C085          C9 49          ; KEY09      CMP #$49
00199 C088          D0 03          BNE KEY10
00200 C08A          4C 1D C2      JMP I01
00201 C08C          C9 4A          ; KEY10      CMP #$4A
00202 C08E          D0 03          BNE KEY11
00203 C090          4C 41 C2      JMP J01
00204 C092          C9 4B          ; KEY11      CMP #$4B
00205 C094          D0 03          BNE KEY12
00206 C096          4C 70 C2      JMP K01
00207 C098          C9 4C          ; KEY12      CMP #$4C
00208 C099          'C'?
00209 C09A          ;NO, KEEP SEARCHING
00210 C09B          ;GO TO 'C' ROUTINE
00211 C09C          'D'?
00212 C09D          ;NO, KEEP SEARCHING
00213 C09E          ;GO TO 'D' ROUTINE
00214 C09F          'E'?
00215 C0A0          ;NO, KEEP SEARCHING
00216 C0A1          ;GO TO 'E' ROUTINE
00217 C0A2          'F'?
00218 C0A3          ;NO, KEEP SEARCHING
00219 C0A4          ;GO TO 'F' ROUTINE
00220 C0A5          'I'?
00221 C0A6          ;NO, KEEP SEARCHING
00222 C0A7          ;GO TO 'I' ROUTINE
00223 C0A8          'J'?
00224 C0A9          ;NO, KEEP SEARCHING
00225 C0AA          ;GO TO 'J' ROUTINE
00226 C0AB          'K'?
00227 C0AC          ;NO, KEEP SEARCHING
00228 C0AD          ;GO TO 'K' ROUTINE
00229 C0AE          'L'?
00230 C0AF          ;NO, KEEP SEARCHING
00231 C0B0          ;GO TO 'L' ROUTINE

```

## C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00216	C09F	D0 03		BNE	KEY13	;NO, KEEP SEARCHING
00217	C0A1	4C 81 C2		JMP	L01	;GO TO 'L' ROUTINE
00218	C0A4		:			
00219	C0A4	C9 4D	KEY13	CMP	#\$4D	; 'M'?
00220	C0A6	D0 03		BNE	KEY14	;NO, KEEP SEARCHING
00221	C0A8	4C BA C2		JMP	M01	;GO TO 'M' ROUTINE
00222	C0AB		:			
00223	C0AB	C9 4E	KEY14	CMP	#\$4E	; 'N'?
00224	C0AD	D0 03		BNE	KEY15	;NO, KEEP SEARCHING
00225	C0AF	4C E1 C2		JMP	N01	;GO TO 'N' ROUTINE
00226	C0B2		:			
00227	C0B2	C9 52	KEY15	CMP	#\$52	; 'R'?
00228	C0B4	D0 03		BNE	KEY16	;NO, KEEP SEARCHING
00229	C0B6	4C FB C2		JMP	R01	;GO TO 'R' ROUTINE
00230	C0B9		:			
00231	C0B9	C9 53	KEY16	CMP	#\$53	; 'S'?
00232	C0BB	D0 03		BNE	KEY17	;NO, KEEP SEARCHING
00233	C0BD	4C 3B C3		JMP	S01	;GO TO 'S' ROUTINE
00234	C0C0		:			
00235	C0C0	C9 54	KEY17	CMP	#\$54	; 'T'?
00236	C0C2	D0 03		BNE	KEY18	;NO, KEEP SEARCHING
00237	C0C4	4C 74 C3		JMP	T01	;GO TO 'T' ROUTINE
00238	C0C7		:			
00239	C0C7	C9 55	KEY18	CMP	#\$55	; 'U'?
00240	C0C9	D0 03		BNE	KEY19	;NO, KEEP SEARCHING
00241	C0CB	4C 89 C3		JMP	U01	;GO TO 'U' ROUTINE



C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00272	COF9	A5 8C	B02	LDA	INDEX1	; RECALL HEIGHT
00273	COFB	C9 08		CMP	#\$08	; LESS THAN 8?
00274	COFF	30 1A		BMI	B03	; YES, SO GO THERE
00275	COFF	E9 08		SBC	#\$08	; TAKE 8 AWAY
00276	C101	85 8C		STA	INDEX1	; STORE IT
00277	C103	A9 A0		LDA	#\$A0	; LOAD A RVS SPACE
00278	C105	91 4B		STA	(USRO), Y	; STORE ON SCREEN
00279	C107	20 2D C1		JSR	FIXCOL	; STORE COLOUR TOO
00280	C10A	18		CLC		
00281	C10B	A5 4B		LDA	USRO	; GET SCREEN LOW BYTE
00282	C10D	E9 27		SBC	#\$27	; BACK ONE SCREEN LINE
00283	C10F	85 4B		STA	USRO	; STORE IT
00284	C111	A5 4C		LDA	USR1	; GET SCREEN HIGH BYTE
00285	C113	E9 00		SBC	#\$00	; ADJUST FOR CARRY
00286	C115	85 4C		STA	USR1	; STORE IT
00287	C117	D0 E0		BNE	B02	; BRANCH ALWAYS
00288	C119		:			
00289	C119	AA	B03	TAX		; GET FRACTION OF 8
00290	C11A	RD 25 C1		LDA	B04, X	; GET FRACTIONAL CHAR
00291	C11D	91 4B		STA	(USRO), Y	; PUT ONTO SCREEN
00292	C11F	20 2D C1		JSR	FIXCOL	
00293	C122	4C 73 00		JMP	CHRGET	; FINISHED - GET NEXT
00294	C125		:			
00295	C125		:			; TABLE OF FRACTIONAL CHARACTERS
00296	C125		:			; TO CONSTRUCT VERTICAL BAR
00297	C125		:			; ACCURATE TO 1/8 OF 1 SCREEN

```

00298 C125
00299 C125
00300 C125 64
00300 C126 6F
00300 C127 79
00300 C128 62
00301 C129 F8
00301 C12A F7
00301 C12B E3
00301 C12C A0
00302 C12D
00303 C12D A9 DC
00304 C12F 45 4C
00305 C131 85 4C
00306 C133 AD 86 02
00307 C136 91 4B
00308 C138
00309 C138 A9 DC
00310 C13A 45 4C
00311 C13C 85 4C
00312 C13E 60
00313 C13F
00314 C13F
00315 C13F
00316 C13F
00317 C13F
00318 C13F
00319 C13F
00320 C13F
00321 C13F 20 9B B7

;LETTER POSITION
;
B04 .BYT $64,$6F,$79,$62

.BYT $F8,$F7,$E3,$A0

;
; FIXCOL LDA #$DC ;CONVERT SCREEN
; EOR USR1 ;RAM TO COLOUR
; STA USR1
; LDA COLOUR
; STA (USRO),Y ;SET THE COLOUR

; SWPCOL LDA #$DC ;AND BACK TO
; EOR USR1 ;SCREEN RAM
; STA USR1
; RTS

;*****
; * C = CLEAR SCREEN SECTION *
;*****
;
; SYNTAX @C X
; X=0-39
;
; C01 JSR GTBYTE ;GET PARAMETER

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00322	C142	A9 04		LDA	#\$04	; SCREEN HIGH BYTE
00323	C144	85 4C		STA	USR1	; STORE IT
00324	C146	A9 27		LDA	#\$27	; SCREEN WIDTH (0-39)
00325	C148	86 4D		STX	USR2	; STORE PARAMETER
00326	C14A	E5 4D		SBC	USR2	
00327	C14C	69 4F		ADC	#\$4F	
00328	C14E	85 4B		STA	USRO	
00329	C150		; C02			
00330	C150	A9 20		LDA	#\$20	; LOAD A SPACE
00331	C152	91 4B		STA	(USR0),Y	
00332	C154	18		CLC		
00333	C155	A5 4B		LDA	USRO	; SCREEN LOW BYTE
00334	C157	69 28		ADC	#\$28	; ADD 1 SCREEN LINE
00335	C159	85 4B		STA	USRO	; STORE IT
00336	C15B	A5 4C		LDA	USR1	; FIX HIGH BYTE
00337	C15D	69 00		ADC	#\$00	
00338	C15F	85 4C		STA	USR1	; STORE IT
00339	C161	C9 07		CMP	#\$07	; FINISHED SCREEN YET?
00340	C163	D0 EB		BNE	C02	; NO, SO GO BACK
00341	C165	A5 4B		LDA	USRO	; CHECK LOW BYTE
00342	C167	C9 20		CMP	#\$20	
00343	C169	30 E5		BMI	C02	; NO, GO BACK
00344	C16B	4C 73 00		JMP	CHRGET	; FINISH OFF GET NEXT
00345	C16E		; *****			
00346	C16E		; * D = SCROLL SCREEN DOWN *			
00347	C16E		; *****			
00348	C16E		; *****			

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00349 C16E      ; SYNTAX @D
00350 C16E      ;
00351 C16E      ; SET UP START AND END OF SCREEN TO BE MOVED
00352 C16E      ; SCREEN START IS AT LEFT END OF FIRST SCREEN LINE
00353 C16E      ; SCREEN END IS SECOND LAST SCREEN LINE
00354 C16E      ; ALL LINES MOVE DOWN ONE LINE
00355 C16E      ; TOP LINE IS FILLED WITH SPACES
00356 C16E      ;
00357 C16E      ;
00358 C16E      A0 28      LDY #28      ; LOAD SCREEN END LOW BYTE
00359 C170      84 8E      STY INDEX3  ; STORE IT
00360 C172      A0 00      LDY #00      ; LOAD LOW BYTE ON NEXT LINE UP
00361 C174      84 8C      STY INDEX1  ; STORE IT
00362 C176      A0 BF      LDY #BF      ; SET INDEX
00363 C178      A2 07      LDX #07      ; OF 2ND LAST LINE
00364 C17A      ;
00365 C17A      86 8D      STX INDEX2  ; STORE IT
00366 C17C      86 8F      STX INDEX4  ; TWICE
00367 C17E      ;
00368 C17E      B1 8C      LDA (INDEX1),Y ; GET CHAR FROM LINE 'L'
00369 C180      91 8E      STA (INDEX3),Y ; MOVE IT TO LINE 'L+1'
00370 C182      20 A2 C1   JSR SWAPC2   ; SET COLOUR RAM
00371 C185      B1 8C      LDA (INDEX1),Y
00372 C187      91 8E      STA (INDEX3),Y
00373 C189      20 A2 C1   JSR SWAPC2
00374 C18C      88        DEY          ; SET SCREEN RAM
00375 C18D      C0 FF      CPY #FF      ; COUNT INDEX DOWN
00376 C18F      D0 ED      BNE D03     ; IS Y < 0?
00377 C191      CA        DEX          ; NO
                                ; COUNT MAIN DOWN

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00378	C192	E0 03		CPX	#\$03	; RUN OFF SCREEN YET?
00379	C194	D0 E4		BNE	D02	; NO
00380	C196	A0 27		LDY	#\$27	; SET COUNTER TO 1 LINE LENGTH
00381	C198	A9 20		LDA	#\$20	; LOAD A SPACE
00382	C19A		;			
00383	C19A	91 8C	D04	STA	(INDEX1),Y	; STORE ON TOP LINE
00384	C19C	88		DEY		; COUNT DOWN INDEX
00385	C19D	10 FB		BPL	D04	; GO BACK FOR MORE
00386	C19F	4C 79 00		JMP	CHRGOT	; YES, SO FINISHED
00387	C1A2		;			
00388	C1A2	A9 DC	SWAPC2	LDA	#\$DC	
00389	C1A4	45 8F		EOR	INDEX4	
00390	C1A6	85 8F		STA	INDEX4	
00391	C1A8		;			
00392	C1A8	A9 DC	SWAPC3	LDA	#\$DC	
00393	C1AA	45 8D		EOR	INDEX2	
00394	C1AC	85 8D		STA	INDEX2	
00395	C1AE	60		RTS		
00396	C1AF		;			
00397	C1AF		;			
00398	C1AF		;			
00399	C1AF		;			
00400	C1AF		;			
00401	C1AF		;			
00402	C1AF		;			
00403	C1AF		;			

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00404 C1AF          A9 02          ; E01
00405 C1AF          A2 08          LDA # $02
00406 C1B1          A8          LDX # $08
00407 C1B3          20 BA FF      TAY
00408 C1B4          20 OD C4    JSR SETLFS
00409 C1B7          98          JSR GETSTR
00410 C1BA          A2 00          TYA
00411 C1BB          A0 CF          LDX # <STRBUF
00412 C1BD          20 BD FF      LDY # >STRBUF
00413 C1BF          20 C0 FF      JSR SETNAM
00414 C1C2          A2 02          JSR OPEN
00415 C1C5          20 C6 FF      LDX # $02
00416 C1C7          C1CA          JSR CHKIN
00417 C1CA          20 CF FF      ; E02
00418 C1CA          20 D2 FF      JSR INCHR
00419 C1CD          24 90          JSR OUTCH
00420 C1D0          70 11          BIT SATUS
00421 C1D2          C9 OD          BVS E04
00422 C1D4          D0 F2          CMP # $0D
00423 C1D6          AD 8D 02        BNE E02
00424 C1D8          29 01          ; E03
00425 C1D8          D0 F9          LDA SFST
00426 C1DB          A5 C5          AND # $01
00427 C1DD          C9 3F          BNE E03
00428 C1DF          D0 E5          LDA LSTX
00429 C1E1          20 CC FF      CMP # $3F
00430 C1E3          A9 02          BNE E02
00431 C1E5          20 CC FF      ; E04
00432 C1E5          C1E8          JSR CLRCHN
00433 C1E8          A9 02          LDA # $02
; SETUP OPEN
; TO DISK
; AS 'OPEN 2,8,2'
; GET FILENAME
; LENGTH OF FN
; POINTER TO
; FILENAME
; OPEN 2,8,2, FN
; OPEN CHANNEL
; GET CHR FROM FILE
; PRINT IT
; EOF ?
; YES
; END OF LINE?
; SHIFT KEY
; NOT CTRL KEY
; LOOP, AND LOOP...
; CHECK RUN/STOP
; NOT PRESSED - GO ON
; BACK TO KEYBOARD

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00434	C1EA	20 C3 FF		JSR	CLOSE	;CLOSE 2 AND EXIT
00435	C1ED	4C 73 00		JMP	CHRGET	
00436	C1F0					
00437	C1F0					*****
00438	C1F0					* F = FILL SCREEN WITH CHARACTER *
00439	C1F0					*****
00440	C1F0					
00441	C1F0					; SYNTAX @F CH,C
00442	C1F0					; CH IS SCREEN POKE OF CHARACTER
00443	C1F0					; C IS SCREEN COLOUR OF CHARACTER
00444	C1F0					
00445	C1F0	20 9B B7	F01	JSR	GTBYTE	;GET CHARACTER REQUIRED
00446	C1F3	86 4B		STX	USRO	;STORE IT
00447	C1F5	20 F1 B7		JSR	COMBYT	
00448	C1F8	8E 86 02		STX	COLOUR	
00449	C1FB	A0 00		LDY	#\$00	;LOAD SCREEN LOW BYTE
00450	C1FD	84 8C		STY	INDEX1	;STORE IT
00451	C1FF	A2 04		LDX	#\$04	;LOAD SCREEN HIGH BYTE
00452	C201	A5 4B		LDA	USRO	;GET CHARACTER BACK
00453	C203					
00454	C203	86 8D	F02	STX	INDEX2	;STORE SCREEN HIGH BYTE
00455	C205					
00456	C205	91 8C	F03	STA	(INDEX1),Y	;STORE CHAR ON SCREEN
00457	C207	20 A8 C1		JSR	SWAPC3	;SET TO COLOUR RAM
00458	C20A	AD 86 02		LDA	COLOUR	
00459	C20D	91 8C		STA	(INDEX1),Y	;STORE COLOUR



C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00490	C239	E8		INX		; INCREMENT HIGH BYTE
00491	C23A	86 4C		STX	USR1	
00492	C23C	E0 08		CPX	#\$08	; FINISHED SCREEN YET?
00493	C23E	D0 ED		BNE	IO3	; NO, GO BACK FOR MORE
00494	C240	60		RTS		; YES, SO FINISHED
00495	C241					;
00496	C241					; *****
00497	C241					; * J = CENTRE STRING ON SCREEN *
00498	C241					; *****
00499	C241					;
00500	C241					; SYNTAX @J "MESSAGE"
00501	C241					;
00502	C241	20 0D C4	J01	JSR	GETSTR	; GET MESSAGE
00503	C244	A9 28		LDA	#\$28	; LINE LENGTH
00504	C246	84 4D		STY	USR2	; STRING LENGTH
00505	C248	E5 4D		SBC	USR2	; 40-LENGTH (CARRY'S SET)
00506	C24A	85 4D		STA	USR2	; STORE IT
00507	C24C					;
00508	C24C	A9 20	J02	LDA	#\$20	; LOAD A SPACE
00509	C24E	20 D2 FF		JSR	OUTCH	; PRINT IT
00510	C251	C6 4D		DEC	USR2	; TAKE 2 AWAY FOR
00511	C253	C6 4D		DEC	USR2	; EVERY SPACE PRINTED
00512	C255	A5 4D		LDA	USR2	; CHECK ITS VALUE
00513	C257	C9 01		CMP	#\$01	; IS IT 1 YET?
00514	C259	10 F1		BPL	J02	; NO, SO PRINT MORE SPACES
00515	C25B	A0 00		LDY	#\$00	; SET INDEX

```

00516 C25D ;
00517 C25D B9 00 CF ; J04 LDA STRBUF, Y ;GET STRING CHAR
00518 C260 F0 06 BEQ J05 ;DONE!
00519 C262 20 D2 FF JSR OUTCH ;PRINT STRING CHAR
00520 C265 C8 INY ;INCREASE POINTER
00521 C266 D0 F5 BNE J04 ;ALWAYS
00522 C268 ;
00523 C268 J05 LDA #$0D ;
00524 C26A 20 D2 FF JSR OUTCH ;
00525 C26D 4C 73 00 JMP CHRGET ;
00526 C270 ;
00527 C270 ;*****
00528 C270 ;* K = KILL REPEAT *
00529 C270 ;*****
00530 C270 ;
00531 C270 ;SYNTAX @K
00532 C270 ;
00533 C270 K01 LDA #<KEY ;IRQ LOW BYTE
00534 C272 8D 14 03 STA CINV1 ;STORE IT
00535 C275 A9 EA LDA #>KEY ;IRQ HIGH BYTE
00536 C277 8D 15 03 STA CINV2 ;STORE IT
00537 C27A ;
00538 C27A K02 LDA #$10 ;
00539 C27C 85 FC STA TEMP1 ;AND OUT AGAIN
00540 C27E 4C 79 00 JMP CHRGT ;
00541 C281 ;
00542 C281 ;*****
00543 C281 ;* L = LOAD SCREEN FROM RAM *
00544 C281 ;*****
00545 C281 ;

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00546	C281					
00547	C281					
00548	C281					
00549	C281					
00550	C281					
00551	C281					
00552	C281					
00553	C281					
00554	C281					
00555	C281					
00556	C281					
00557	C281					
00558	C281					
00559	C281	A2 04	L01	LDX	#\$04	; SCREEN HIGH BYTE
00560	C283	A0 00		LDY	#\$00	; SCREEN LOW BYTE
00561	C285	84 8C		STY	INDEX1	; STORE
00562	C287	84 8E		STY	INDEX3	; LOW BYTES
00563	C289	A9 C3		LDA	#\$C3	; START ADDRESS-2 BLOCKS
00564	C28B	85 8F		STA	INDEX4	
00565	C28D					
00566	C28D	86 8D		STX	INDEX2	; STORE HIGH BYTE
00567	C28F	E6 8F		INC	INDEX4	; SKIP LAST
00568	C291	E6 8F		INC	INDEX4	; TWO BLOCKS
00569	C293	A9 00		LDA	#\$00	; LOAD SCREEN STORE LOW BYTE
00570	C295	85 8E		STA	INDEX3	; STORE IT
00571	C297					

```

00572 C297 ;(INDEX3) POINTS TO $C500
00573 C297 ;
00574 C297 ;NOW TO TRANSFER CONTENTS TO SCREEN
00575 C297 ;
00576 C297 L03 LDA (INDEX3),Y ;GET A STORED CHAR
00577 C299 91 8C STA (INDEX1),Y ;PUT ONTO SCREEN
00578 C29B 20 A8 C1 JSR SWAPC3 ;POINT TO COLOUR RAM
00579 C29E E6 8F INC INDEX4 ;AND NEXT BLOCK
00580 C2A0 B1 8E LDA (INDEX3),Y ;SAVE COLOUR
00581 C2A2 91 8C STA (INDEX1),Y ;RAM AND
00582 C2A4 20 A8 C1 JSR SWAPC3 ;BACK TO SCREEN
00583 C2A7 C6 8F DEC INDEX4 ;AND ORIGINAL BLOCK
00584 C2A9 C8 INY ;INCREASE COUNT
00585 C2AA D0 EB BNE L03 ;BLOCK NOT MOVED
00586 C2AC E8 INX ;ADD 1 TO HIGH BYTE
00587 C2AD E0 08 CPX #$08 ;FINISHED SCREEN YET?
00588 C2AF D0 DC BNE L02 ;NO, GO BACK FOR MORE
00589 C2B1 AD FE CF LDA COLTMP+1
00590 C2B4 8D 21 D0 STA SCCOL
00591 C2B7 4C 79 00 JMP CHRGT
00592 C2BA ;
00593 C2BA ;*****
00594 C2BA ;* M = REPLACE CHARACTER *
00595 C2BA ;*****
00596 C2BA ;
00597 C2BA ;SYNTAX @M A,B
00598 C2BA ; A IS PRESENT CHAR ON SCREEN
00599 C2BA ; B IS REPLACEMENT CHAR ON SCREEN
00600 C2BA ;(CHARACTERS ARE SCREEN POKE VALUES)
00601 C2BA ;

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C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00602	C2BA	20 9B B7	M01	JSR	GTBYTE	; GET OLD CHAR
00603	C2BD	86 4C		STX	USR1	; STORE IT
00604	C2BF	20 F1 B7		JSR	COMBYT	; GET NEW CHAR
00605	C2C2	86 4B		STX	USRO	; STORE IT
00606	C2C4	A0 00		LDY	#\$00	; SCREEN LOW BYTE
00607	C2C6	84 8C		STY	INDEX1	; STORE IT
00608	C2C8	A2 04		LDX	#\$04	; SCREEN HIGH BYTE
00609	C2CA	A5 4B		LDA	USRO	
00610	C2CC		;			
00611	C2CC	86 8D	M02	STX	INDEX2	; STORE HIGH BYTE
00612	C2CE		;			
00613	C2CE	B1 8C	M03	LDA	(INDEX1),Y	; GET CHAR FROM SCREEN
00614	C2D0	C5 4C		CMP	USR1	; WANT THIS ONE?
00615	C2D2	D0 04		BNE	M04	; NO, SO BRANCH
00616	C2D4	A5 4B		LDA	USRO	; GET NEW CHAR
00617	C2D6	91 8C		STA	(INDEX1),Y	; STORE IT
00618	C2D8		;			
00619	C2D8	C8	M04	INY		; INCREASE COUNTER
00620	C2D9	D0 F3		BNE	M03	
00621	C2DB	E8		INX		
00622	C2DC	E0 08		CPX	#\$08	; FINISHED SCREEN YET?
00623	C2DE	D0 EC		BNE	M02	; NO
00624	C2E0	60		RTS		; YES, SO FINISHED
00625	C2E1		;			
00626	C2E1		;			*****
00627	C2E1		;			* N = DEC TO HEX CONVERSION *
00628	C2E1		;			*****

```

00629 C2E1
00630 C2E1
00631 C2E1
00632 C2E1
00633 C2E1
00634 C2E4
00635 C2E7
00636 C2E8
00637 C2EA
00638 C2ED
00639 C2EE
00640 C2F1
00641 C2F2
00642 C2F5
00643 C2F7
00644 C2FA
00645 C2FB
00646 C2FB
00647 C2FB
00648 C2FB
00649 C2FB
00650 C2FB
00651 C2FB
00652 C2FB
00653 C2FC
00654 C2FE
00655 C301
00656 C303
00657 C306

20 9E AD
20 F7 B7
48
A9 24
A9 24 FF
68
20 3A C4
98
20 3A C4
A9 0D
20 D2 FF
60

; SYNTAX @N X
; X = 0 -> 65535
;
; NO1 JSR FRMNUM
; JSR GETADR
; PHA
; LDA #$24
; JSR OUTCH
; PLA
; JSR WROB
; TYA
; JSR WROB
; LDA #$0D
; JSR OUTCH
; RTS

; *****
; * R = REPEAT KEY *
; *****
; SYNTAX @R
;
; RO1 SEI
; LDA #<R02
; STA CINV1
; LDA #>R02
; STA CINV2
; LDA #$01

20 9E AD ; EVALUATE NUMERIC EXPR.
20 F7 B7 ; CONVERT TO 2-BYTE INTEGER
48 ; STORE IT
A9 24 ; '$'
A9 24 FF ; PRINT IT
68 ; GET IT BACK
20 3A C4 ; OUTPUT 1 BYTE
98 ; GET SECOND BYTE
20 3A C4 ; OUTPUT 1 BYTE
A9 0D ; FINISHED
20 D2 FF
60

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00658	C308	85 4E		STA	REPD L	
00659	C30A	58		CLI		
00660	C30B	4C 7A C2		JMP	K02	; TO FINISH OFF
00661	C30E		;			
00662	C30E	A5 C5	R02	LDA	LSTX	; GET A KEY
00663	C310	C5 FE		CMP	TEMP3	; SAME AS LAST KEY?
00664	C312	F0 09		BEQ	R04	; YES
00665	C314	85 FE		STA	TEMP3	; STORE IT
00666	C316	A9 10		LDA	#\$10	; SET DELAY
00667	C318	85 FD		STA	TEMP2	; STORE IT
00668	C31A		;			
00669	C31A	4C 31 EA	R03	JMP	KEY	; FINISHED JOB
00670	C31D		;			
00671	C31D	C9 FF	R04	CMP	#\$FF	
00672	C31F	F0 F9		BEQ	R03	
00673	C321	A5 FD		LDA	TEMP2	; RECALL DELAY
00674	C323	F0 04		BEQ	R05	; REACHED 0
00675	C325	C6 FD		DEC	TEMP2	; DECREASE COUNTER
00676	C327	D0 F1		BNE	R03	
00677	C329		;			
00678	C329	C6 4E	R05	DEC	REPD L	
00679	C32B	D0 ED		BNE	R03	
00680	C32D	A9 04		LDA	#\$04	
00681	C32F	85 4E		STA	REPD L	
00682	C331	A9 40		LDA	#\$40	; NIL KEY
00683	C333	85 C5		STA	LSTX	; STORE IN KEY REGISTER

```

00684 C335 A9 02 LDA #S02
00685 C337 85 CD STA BLNCT
00686 C339 D0 DF BNE R03
00687 C33B
00688 C33B
00689 C33B
00690 C33B
00691 C33B
00692 C33B
00693 C33B
00694 C33B
00695 C33B
00696 C33B
00697 C33B
00698 C33B
00699 C33B
00700 C33B
00701 C33E AD 21 D0
00702 C341 8D FE CF
00703 C343 A2 04
00704 C345 A0 00
00705 C347 84 8C
00706 C349 84 8E
00707 C34B A9 C3
00708 C34D 85 8F
00709 C34D 86 8D
00710 C34F E6 8F
00711 C351 E6 8F
00712 C353 A9 00
00713 C355 85 8E

;
; *****
; * S = SAVE SCREEN IN RAM *
; *****
;
; SYNTAX @S
;
; THIS ROUTINE TAKES THE SCREEN CONTENTS
; AND MOVES THEM TO A LOCATION IN RAM
; STARTING AT $C500, FROM WHERE THEY
; CAN BE RECALLED WITH '@L' OR SWAPPED
; WITH '@W'.
;
; S01 LDA SCCOL ;REMEMBER SCREEN
STA COLTMP+1 ;COLOUR
LDX #S04 ;SCREEN HIGH BYTE
LDY #S00 ;SCREEN LOW BYTE
STY INDEX1 ;STORE
STY INDEX3 ;LOW BYTES
LDA #C3 ;START ADDRESS-2 BLOCKS
STA INDEX4
;
; S02 STX INDEX2 ;STORF. HIGH BYTE
INC INDEX4 ;SKIP LAST
INC INDEX4 ;TWO BLOCKS
LDA #S00 ;LOAD SCREEN STORE LOW BYTE
STA INDEX3 ;STORE IT

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00714	C357					
00715	C357		;			
00716	C357		;	(INDEX3)	POINTS TO \$C500	
00717	C357		;			
00718	C357		;			
00719	C357		;			
00720	C359	B1 8C	S03	LDA	(INDEX1),Y ;GET A STORED CHAR	
00721	C35B	91 8E		STA	(INDEX3),Y ;PUT INTO RAM	
00722	C35E	20 A8	C1	JSR	SWAPC3 ;POINT TO COLOUR RAM	
00723	C360	E6 8F		INC	INDEX4 ;AND NEXT BLOCK	
00724	C362	B1 8C		LDA	(INDEX1),Y ;SAVE COLOUR	
00725	C364	91 8E	C1	STA	(INDEX3),Y ;RAM AND	
00726	C367	20 A8	C1	JSR	SWAPC3 ;BACK TO SCREEN	
00727	C369	C6 8F		DEC	INDEX4 ;AND ORIGINAL BLOCK	
00728	C36A	D0 EB		INY	INX ;INCREASE COUNT	
00729	C36C	E8		BNE	S03 ;BLOCK NOT MOVED	
00730	C36D	F0 08		INX	INX ;ADD 1 TO HIGH BYTE	
00731	C36F	D0 DC		CPX	#\$08 ;FINISHED SCREEN YET?	
00732	C371	4C 79	00	BNE	S02 ;NO, GO BACK FOR MORE	
00733	C374			JMP	CHRGOT	
00734	C374		;			*****
00735	C374		;			* T = SET CURSOR LOCATION *
00736	C374		;			*****
00737	C374		;			
00738	C374		;			;
00739	C374		;			;
						SYNTAX @T X,Y
						; X IS HORIZONTAL POSITION (COLUMN)

```

00740 C374 ; Y IS VERTICAL POSITION (LINE)
00741 C374 ;
00742 C374 ; T01
00743 C377 JSR GTBYTE ;GET X
00744 C379 STX USRO ;STORE IT
00745 C37C JSR COMBYT ;GET Y
00746 C37E STX USR1 ;STORE IT
00747 C380 LDX USRO ;GET X
00748 C382 LDY USR1 ;GET Y
00749 C383 CLC ;SETUP FOR KERNAL
00750 C386 JSR PLOT ;SET CURSOR & EXIT
00751 C389 JMP CHRGET
;
; *****
; * U = MOVE SCREEN UP 1 LINE *
; *****
;
; SYNTAX @U
;
; ROUTINE STARTS WITH FIRST CHAR ON
; SECOND LINE OF SCREEN, MOVING ALL
; LINES UP ONE. THE BOTTOM LINE IS THEN
; FILLED WITH SPACES.
;
; U01 LDX #$00 ;LINE 1 SCREEN LOW BYTE
00763 C389 STX INDEX3 ;STORE IT
00764 C38B LDX #$28 ;LINE 2 SCREEN LOW BYTE
00765 C38D STX INDEX1 ;STORE IT
00766 C38F LDX #$04 ;SCREEN HIGH BYTE
00767 C391
00768 C393 ;
00769 C393 U02 ; STORE IT

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00770	C395	86 8D		STX	INDEX2	;.. IN BOTH PLACES
00771	C397	A0 00		LDY	#\$00	;SET COUNTER
00772	C399					
00773	C399	B1 8C	; U03	LDA	(INDEX1),Y	;GET CHAR
00774	C39B	91 8E		STA	(INDEX3),Y	;MOVE UP 1 LINE
00775	C39D	20 A2 C1		JSR	SWAPC2	
00776	C3A0	B1 8C		LDA	(INDEX1),Y	
00777	C3A2	91 8E		STA	(INDEX3),Y	
00778	C3A4	20 A2 C1		JSR	SWAPC2	
00779	C3A7	C8		INY		
00780	C3A8	D0 EF		BNE	U03	
00781	C3AA	E8		INX		
00782	C3AB	E0 08		CPX	#\$08	;FINISHED SCREEN YET
00783	C3AD	D0 E4		BNE	U02	;NO, SO LOOP BACK
00784	C3AF	A9 20		LDA	#\$20	;YES, SO LOAD A SPACE
00785	C3B1	A2 28		LDX	#\$28	;SET A COUNTER
00786	C3B3					
00787	C3B3	9D C7 07	; U04	STA	\$07C7,X	;ON LAST SCREEN LINE
00788	C3B6	CA		DEX		;COUNT DOWN
00789	C3B7	D0 FA		BNE	U04	;NOT DONE YET
00790	C3B9	4C 79 00		JMP	CHRGOT	
00791	C3BC					
00792	C3BC		; *****			
00793	C3BC		; * W = SWAP SCREEN WITH RAM STORE *			
00794	C3BC		; *****			
00795	C3BC		; *****			

```

00796 C3BC
00797 C3BC
00798 C3BC
00799 C3BC
00800 C3BC
00801 C3BC
00802 C3BC
00803 C3BC
00804 C3BC
00805 C3BC
00806 C3BF
00807 C3C2
00808 C3C4
00809 C3C6
00810 C3C8
00811 C3CA
00812 C3CC
00813 C3CE
00814 C3CE
00815 C3D0
00816 C3D2
00817 C3D4
00818 C3D4
00819 C3D6
00820 C3D8
00821 C3DA
00822 C3DC
00823 C3DE
00824 C3E0
00825 C3E3

AD 21 D0
8D FD CF
A2 04
A0 00
84 8C
84 8E
A9 C3
85 8F
86 8D
E6 8F
E6 8F
B1 8E
85 4B
B1 8C
91 8E
A5 4B
91 8C
20 A8 C1
E6 8F

;SYNTAX @W
;
;THIS ROUTINE DOES A BYTE BY BYTE SWAP
;OF SCREEN CONTENTS WITH THE CONTENTS
;OF RAM - MAYBE ANOTHER SCREEN STORED
;THERE. THE SWAP IS A 3-MOVE S/RTNE
;SCREEN BYTE TO TEMP STORE,
;RAM TO SCREEN, THEN TEMP TO RAM.
;
; W01 LDA SCCOL
STA COLTMP
LDX #$04
LDY #$00
STY INDEX1
STY INDEX3
LDA #$C3
STA INDEX4
;
; W02 STX INDEX2
INC INDEX4
INC INDEX4
;
; W03 LDA (INDEX3),Y
STA USRO
LDA (INDEX1),Y
STA (INDEX3),Y
LDA USRO
STA (INDEX1),Y
JSR SWAPC3
INC INDEX4
; GET FROM RAM
; STASH IT SAFELY
; GET FROM SCREEN
; PUT IN RAM
; PUT ON SCREEN
; POINT TO COLOUR RAM
; POINT TO COPY

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00826	C3E5	B1 8E		LDA	(INDEX3),Y	;GET FROM RAM COPY
00827	C3E7	85 4B		STA	USRO	;SAFE FOR SWAP
00828	C3E9	B1 8C		LDA	(INDEX1),Y	;GET THE REAL THING
00829	C3EB	91 8E		STA	(INDEX3),Y	;COPY TO RAM
00830	C3ED	A5 4B		LDA	USRO	
00831	C3EF	91 8C		STA	(INDEX1),Y	;PUT IN COLOUR RAM
00832	C3F1	C6 8F		DEC	INDEX4	;POINT BACK TO SCREEN COPY
00833	C3F3	20 A8 C1		JSR	SWAPC3	
00834	C3F6	C8		INY		;BACK TO SCREEN
00835	C3F7	D0 DB		BNE	W03	
00836	C3F9	E8		INX		
00837	C3FA	E0 08		CPY	#\$08	
00838	C3FC	D0 D0		BNE	W02	
00839	C3FE	AD FE CF		LDA	COLTMP+1	
00840	C401	8D 21 D0		STA	SCCOL	
00841	C404	AD FD CF		LDA	COLTMP	
00842	C407	8D FE CF		STA	COLTMP+1	
00843	C40A	4C 79 00		JMP	CHRGOT	
00844	C40D					
00845	C40D					
00846	C40D					
00847	C40D					
00848	C40D					
00849	C40D	20 73 00		GETSTR	JSR CHRGET	;SCAN FOR
00850	C410	C9 22		CMP	#\$22	;FIRST QUOTE
00851	C412	F0 03		BEQ	STRTOK	

```

00852 C414 4C 08 AF JUNK JMP SNERR ;UNREADABLE
00853 C417 A0 01 ;STR TOK LDY #01
00854 C417 A0 01 ;
00855 C419 B1 7A ;STASH LDA (CHRG7),Y ;GET STRING
00856 C419 B1 7A ;
00857 C41B F0 F7 ;STR TOK LDY #01 ;END OF LINE ?
00858 C41D C9 22 ;STASH LDA (CHRG7),Y ;END QUOTE ?
00859 C41F F0 06 ;STR TOK LDY #01 ;YES - FINISH UP
00860 C421 99 FF CE ;STASH LDA (CHRG7),Y ;STORE IN BUFFER
00861 C424 C8 ;STR TOK LDY #01 ;
00862 C425 D0 F2 ;STR TOK LDY #01 ;ALWAYS BACK FOR MORE
00863 C427 18 ;STR TOK LDY #01 ;
00864 C428 A9 00 ;ENDSTR CLC ;END-OF-STRING POINTER
00865 C428 A9 00 ;STR TOK LDY #01 ;
00866 C42A 99 FF CE ;STR TOK LDY #01 ;
00867 C42D 98 ;STR TOK LDY #01 ;
00868 C42E 65 7A ;STR TOK LDY #01 ;SET CHRGOT
00869 C430 85 7A ;STR TOK LDY #01 ;TO END OF
00870 C432 A9 00 ;STR TOK LDY #01 ;
00871 C434 65 7B ;STR TOK LDY #01 ;LINE
00872 C436 85 7B ;STR TOK LDY #01 ;LENGTH IS ONE LESS
00873 C438 88 ;STR TOK LDY #01 ;
00874 C439 60 ;STR TOK LDY #01 ;
00875 C43A ;STR TOK LDY #01 ;
00876 C43A ;STR TOK LDY #01 ;
00877 C43A ;STR TOK LDY #01 ;
00878 C43A ;STR TOK LDY #01 ;
00879 C43A ;STR TOK LDY #01 ;
00880 C43A 48 ;STR TOK LDY #01 ;SAVE LOW BITS
00881 C43B 4A ;STR TOK LDY #01 ;GET UPPER BITS

```

C64 DEMONSTRATION

LINE#	LOCN	M/CODE	LABEL	OPC	OPERAND	COMMENT
00882	C43C	4A		LSR	A	
00883	C43D	4A		LSR	A	
00884	C43E	4A		LSR	A	
00885	C43F	C9 0A		COMP	#\$0A	; AN ALPHA-HEX CHR ?
00886	C441	90 02		BCC	NUM1	; NO, JUST A NUMBER
00887	C443	69 06		ADC	#\$06	; ADD 7, AS CARRY'S SET
00888	C445		;			
00889	C445	69 30	NUM1	ADC	#\$30	; BUT NOW IT'S CLEAR
00890	C447	20 D2 FF		JSR	OUTCH	; PRINT IT
00891	C44A	68		PLA		; GET LOW NYBBLE
00892	C44B	29 0F		AND	#\$0F	; MASK OFF HIGH BITS
00893	C44D	C9 0A		COMP	#\$0A	; CHECK FOR ALPHA AGAIN
00894	C44F	90 02		BCC	NUM2	
00895	C451	69 06		ADC	#\$06	
00896	C453		;			
00897	C453	69 30	NUM2	ADC	#\$30	
00898	C455	4C D2 FF		JMP	OUTCH	; PRINT & EXIT
00899	C458		;			
00900	C458		;			; ---SET ASIDE A STRING BUFFER---
00901	C458		;			
00902	C458		STRBUF	=\$CF00		
00903	C458		;			
00904	C458					.END

ERRORS = 00000



LABEL INDEX

LABEL	LOCN								
R04	C31D								004E
S01	C33B								0090
SCCOL	D021								FFBD
SFST	028D								CF00
STR TOK	C417								C138
T01	C374								00FE
U01	C389								C3B3
USR0	004B								C3BC
W02	C3CE								C00D
R05	C329								
S02	C34D								
SCREEN	0400								
SNERR	AF08								
SWAPC2	C1A2								
TEMP1	00FC								
U02	C393								
USR1	004C								
W03	C3D4								
READY									
S03									
SETLFS									
STASH									
SWAPC3									
TEMP2									
U03									
USR2									
WROB									
A474									
C357									
FFBA									
C419									
C1A8									
00FD									
C399									
004D									
C43A									
REPDL									
SATUS									
SETNAM									
STRBUF									
SWPCOL									
TEMP3									
U04									
W01									
Z01									

END OF ASSEMBLY

# COMMODORE 64 MACHINE LANGUAGE TUTORIAL

## WHAT COMES AFTER BASIC:

BASIC has become the universal computer language for home and personal computers. To the computerist it soon becomes apparent that speed is BASIC's main drawback.

That's where MACHINE LANGUAGE (ML) comes in. By removing the intermediate interpretation of step BASIC, the computer can run hundreds, if not thousands of times faster. For tasks like sorting, searching and some graphics, ML is the fastest (and probably the only satisfactory) way of programming. All tasks that require a large number of repetitions of whatever type – that is where ML will win hands down. You will also find that judicious use of ML permits larger or more complex programs to reside in the limited memory space of a micro computer.

## THE AUTHOR:

Paul Blair (DipCE, MIE Aust.), is a professional engineer in the field of national road policy.

Computing interest started in the sixties when he commenced writing structural design programs in FORTRAN IV. In 1979, he moved into writing custom software. Involvement grew as did international contacts which include Canada, U.S.A., England, New Zealand and the Netherlands. Paul has had items published in most of these countries and is a regular contributor to the Australian Commodore Magazine. He is a strong supporter of computer user groups and is constantly called upon for demonstrations and tutorials in schools and other institutions.

**A second or follow up volume to this book is already being formed.**

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